

Ieke Moerdijk - Dendroidal Topology II

- coloured operad in Sets, in Top.

- for any finite rooted tree T we have introduced the operad $\Omega(T)$, "freely" generated by T .

Today, we will discuss a category of presheaves

$$\text{Sets}^{\Omega^{\text{op}}}$$

on a small category Ω , made up from these trees.

Definition of Ω : objects are finite rooted trees T
arrows are maps of coloured operads $\Omega(S) \rightarrow \Omega(T)$

What does this category look like?

Note that there is a fully faithful embedding

$$\Delta \xrightarrow{i} \Omega$$

from the category Δ of finite linear orders

$$[n] = \{0, \dots, n\} \quad (n \geq 0)$$

$$i[n] = \begin{array}{c} | \\ \bullet \\ \vdots \\ \bullet \\ | \\ 0 \end{array}$$

2) for a corolla

$$C_n =$$



there are external faces



$$i = 0, \dots, k$$

Proposition: Every arrow $S \rightarrow T$ in Ω

can be written essentially uniquely as a composition

$$S \xrightarrow{\sigma} S' \xrightarrow[\alpha]{\sim} T' \xrightarrow{\partial} T$$

where α is an iso, σ is a composition of degeneracies

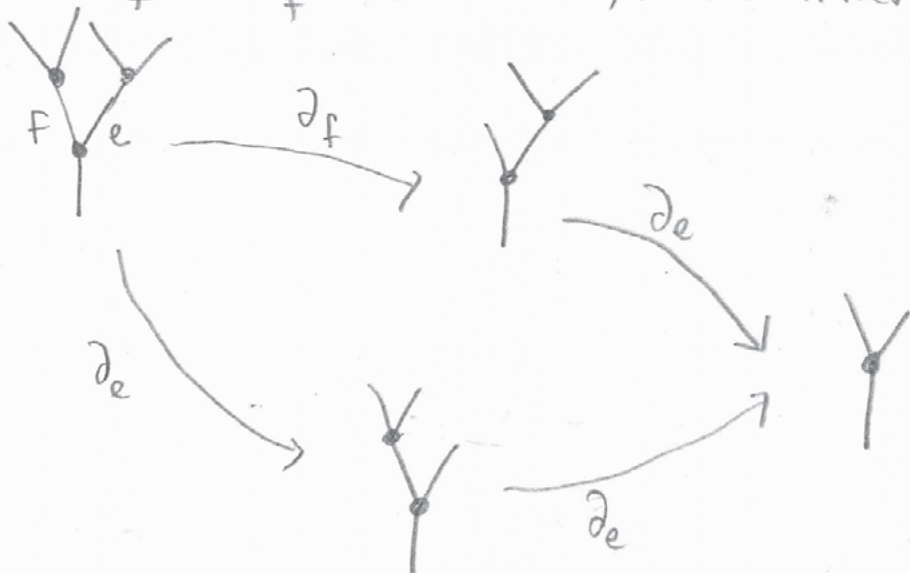
∂ is a composition of faces.

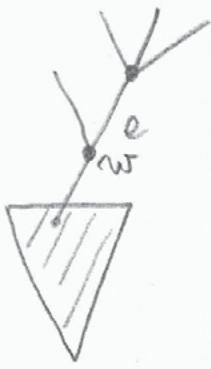
So in this sense Ω is generated by faces and degeneracies as well as symmetries

(in Δ we only have faces & degeneracies).

One could (but I will not) write a list of relations similar to the simplicial identities, eg

$$\partial_e \partial_f = \partial_f \partial_e \text{ if } e, f \text{ are inner edges in a tr.}$$





$$\partial_w \partial_v = \partial_{w \circ_e v} \partial_e$$

if e relates an external vertex v to a vertex w so that e is the only input edge of w which is internal.

Proposition. the classifying space of Ω is contractible.

Proof: let $\bar{\Omega}$ be the full subcategory of "closed" trees, ie trees without leaves. The inclusion has a left adjoint

$$\Omega \begin{array}{c} \xrightarrow{d} \\ \xleftarrow{\quad} \end{array} \bar{\Omega}, \quad d\left(\begin{array}{c} \diagup \quad \diagdown \\ \cdot \quad \cdot \\ \diagdown \quad \diagup \\ \cdot \quad \cdot \\ \diagup \quad \diagdown \\ \cdot \quad \cdot \end{array}\right) = \begin{array}{c} \cdot \quad \cdot \quad \cdot \quad \cdot \\ \diagdown \quad \diagup \\ \cdot \\ \diagdown \quad \diagup \\ \cdot \quad \cdot \end{array}$$

Now on $\bar{\Omega}$ we have the following functors

$$\begin{array}{ccc} & \lambda & \\ \bar{\Omega} & \begin{array}{c} \downarrow d \\ \xrightarrow{id} \\ \uparrow \beta \end{array} & \bar{\Omega} \\ & \text{constant } \uparrow & \end{array} \quad \lambda(T) = \begin{array}{c} \text{shaded triangle} \\ \cdot \\ \uparrow v_T \text{ new root vertex} \end{array}$$

$d: T \hookrightarrow \lambda(T)$ is the obvious embedding

$\beta: \uparrow \hookrightarrow \lambda(T)$ sends the vertex to the new v_T in $\lambda(T)$.

Remark: λ is a functor on $\overline{\Omega}$ but not on Ω .

Presheaves on Ω

Definition: the category of presheaves on Ω will be called that of dendroidal sets, and denoted $d\text{Sets} = \text{Sets}^{\Omega^{\text{op}}}$.

Let's see a few examples:

① every simplicial set can be viewed as a dendroidal set

since $\Delta \xrightarrow{i} \Omega$ induces adjoint functors

$$\begin{array}{ccc} & i_! & \\ & \curvearrowright & \\ s\text{Sets} & \xleftarrow{i^*} & d\text{Sets} \\ & \curvearrowleft & \\ & i_* & \end{array}$$

Since i is full and faithful, so are $i_!$ and i_* .

$i_!(M) =$ "extension by zero".

$$\text{i.e. } i_!(M)_T = \begin{cases} M_n & \text{if } T = i[n] \\ \emptyset & \text{otherwise} \end{cases}$$

In fact, $s\text{Sets}$ is an open subtopos of $d\text{Sets}$.

There is an isomorphism

$$\Omega/1 = \Delta$$

which induces an equivalence of categories

$$d\text{Sets} / \Omega(-, 1) = s\text{Sets}.$$

Under this iso,

$i_!$ corresponds to "forget"

i^* corresponds to "pullback"

Note that $\Omega(-, 1) = i_!(\Delta[0])$.

② every object T in Ω defines a representable presheaf

$$\Omega[T] := \Omega(-, T)$$

note $i_!(\Delta[n]) = \Omega[i[n]]$

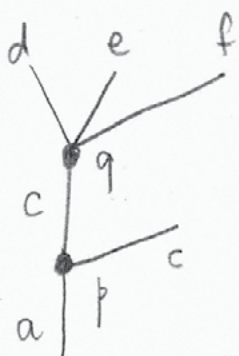
③ every coloured operad P in Sets defines a dendroidal set called its nerve.

$$N(P)_T = \text{Hom}(\Omega(T), P).$$

Remark: an element of $N(P)_T$ looks like a labelling of T :

- its edges by colours of P

- its vertices by operations of P .



More formally, an element of $N(P)_T$ is an equivalence class given by a planar structure on T and such a labelling.

- ④ Every coloured operad P in Top
 (the operations form spaces but there is still a set of colours)

defines a dendroidal set $\mathcal{D}(P)$ defined as

$$\mathcal{D}(P)_T = \text{Hom}(\underbrace{W\Omega(T)}_{W(T)}, P)$$

Recall from Lecture 1 that $\Omega(T)$ is an operad where colours are the edges of T and with exactly one operation $a_1, \dots, a_n \rightarrow b$ if there is a subtree S of T with leaves a_1, \dots, a_n and root b .



Hence elements of $W(T)(a_1, \dots, a_n; b)$ are such subtrees S with lengths

$$\theta \in [0, 1]$$

assigned to the internal edges of S

and these are composed by taking the union as subtrees of T and by assigning length 1 to the edges connecting these subtrees.

$\mathcal{D}(P)$ is called the homotopy coherent nerve of P

this defines a functor

$$\mathcal{D}: (\text{top. operads}) \longrightarrow \text{dSets}$$

It has a left adjoint:

$$W: \text{dSets} \longrightarrow (\text{top. operads})$$

$$W(\Omega[T]) = W(T)$$

W preserves colimits

W is the left Kan extension of the BV resolution applied to the operads $\Omega(T)$ given by representables $\Omega[T]$.

In fact, for an operad P in Sets

$$WN(P) = W(P).$$

$$\begin{array}{ccc} & (\text{Top. operads}) & \\ W \nearrow & & \nwarrow W \\ \Omega & \xrightarrow{\Omega[-]} & \text{dSets} \end{array}$$

The ordinary nerve

$$N: (\text{operads in sets}) \longrightarrow \text{dSets}$$

also has a left adjoint

$$\tau: \text{dSets} \longrightarrow \text{operads}$$

completely determined on representables by

$$\tau(\Omega[T]) = \Omega(\tau).$$

$$\begin{array}{ccc} \text{dSets} & \begin{array}{c} \xrightarrow{\tau} \\ \xleftarrow{N} \end{array} & (\text{operads in sets}) \\ \begin{array}{c} \uparrow i_! \\ \downarrow i^* \end{array} & & \begin{array}{c} \uparrow j_! \\ \downarrow j^* \end{array} \\ \text{sSets} & \begin{array}{c} \xrightarrow{\tau} \\ \xleftarrow{N} \end{array} & (\text{categories}) \end{array}$$

$$\begin{array}{ccc} (\text{top. operads}) & \begin{array}{c} \xleftarrow{W} \\ \xrightarrow{D} \end{array} & \text{dSets} \\ \begin{array}{c} \uparrow \\ \downarrow \end{array} & & \begin{array}{c} \uparrow i_! \\ \downarrow i^* \end{array} \\ (\text{top cats}) & \begin{array}{c} \xleftarrow{W} \\ \xrightarrow{\quad} \end{array} & \text{sSets} \end{array}$$

classical
homotopy
coherent nerve

$$\begin{array}{ccc} & (\text{topological categories}) & \\ W \nearrow & & \nwarrow W \\ \Delta & \xrightarrow{\Delta[-]} & \text{sSets} \end{array}$$

Precise description of $W(\mathbb{C})$ for a category \mathbb{C} ,
 the Boardman-Vogt resolution of \mathbb{C} viewed as an operad $P_{\mathbb{C}}$.
 It is a topological category (cf. lecture 4)

Its objects are the same as of \mathbb{C}

An arrow $c \rightarrow d$ is an equivalence class
 of sequences:

$$c = c_n \xrightarrow{f_n} c_{n-1} \xrightarrow{\theta_{n-1} f_{n-1}} c_{n-2} \xrightarrow{\theta_{n-2}} \dots \xrightarrow{f_2} c_1 \xrightarrow{\theta_1 f_1} c_0 = d$$

$\theta_i \in [0, 1]$ are "waiting times".

If $\theta_i = 0$, the sequence represents the same arrow

as

$$c = c_n \xrightarrow{f_n} c_{n-1} \xrightarrow{\theta_{n-1}} \dots \xrightarrow{\theta_{i+1} f_i \circ f_{i-1}} c_{i-1} \xrightarrow{\theta_{i-1}} \dots \rightarrow c_0$$

Composition is defined by writing length 1
 on the new intermediate object.

Eq: if $\mathbb{C} = [n] = n \rightarrow n-1 \rightarrow \dots \rightarrow 0$

an arrow $i \rightarrow j$ can always be represented
 by a full sequence

$$i \rightarrow i-1 \rightarrow \dots \rightarrow j+1 \rightarrow j$$

as long as we put 0's.

$$(W[n])(i, j) = [0, 1]^{j-1-i}$$

This relates to the classical description of the homotopy coherent nerve of a topological category \mathcal{C} as

$$\mathrm{HcN}(\mathcal{C})_n = \mathrm{Hom}(W[n], \mathcal{C}).$$

Tomorrow's lecture:

there is a symmetric monoidal closed structure $\otimes, \underline{\mathrm{Hom}}$ on $d\mathrm{Sets}$,

extending the cartesian structure on $s\mathrm{Sets}$

$$s\mathrm{Sets} \xrightarrow{i_!} d\mathrm{Sets}$$

$$i_!(M \times N) = i_!(M) \otimes i_!(N)$$

its unit is $\eta = \Omega[1]$.