On the structure of classical realizability models of ZF

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Abstract. In [4, 5, 6], we have introduced the technique of classical realizability, which permits to extend the Curry-Howard correspondence between proofs and programs, to Zermelo-Fraenkel set theory. The models of ZF we obtain in this way, are called realizability models; this technique is an extension of the method of forcing, in which the ordered sets (sets of conditions) are replaced with more complex first order structures called realizability algebras.

We show here that every realizability model \( \mathcal{N} \) of ZF contains a transitive submodel, which has the same ordinals as \( \mathcal{N} \), and which is an elementary extension of the ground model. It follows that the constructible universe of a realizability model is an elementary extension of the constructible universe of the ground model.

We obtain this result by showing the existence of an ultrafilter on the characteristic Boolean algebra \( \mathfrak{A} \) of the realizability model, which is defined in [5, 6].

Introduction

We use here the basic notions and notations of the theory of classical realizability, which was developed in [4, 5, 6].

We consider a model \( \mathcal{M} \) of ZF + V = L, which we call the ground model\(^1\) and, in \( \mathcal{M} \), a realizability algebra \( \mathcal{A} = (\Lambda, \Pi, \Lambda \star \Pi, QP, \bot) \).

\( \Lambda \) is the set of terms, \( \Pi \) is the set of stacks, \( \Lambda \star \Pi \) is the set of processes, \( QP \subset \Lambda \) is the set of proof-like terms, and \( \bot \) is a distinguished subset of \( \Lambda \star \Pi \).

They satisfy the axioms of realizability algebra, which are given in [4] or [6].

In the model \( \mathcal{M} \), we use the language of ZF with the binary relation symbols \( \in, \subset \) and function symbols, which we shall define when needed, by means of formulas of ZF.

We can now build (see [4]) the realizability model \( \mathcal{N} \), which has the same set of individuals as \( \mathcal{M} \), the truth value set of which is \( \mathcal{P}(\Pi) \), endowed with a suitable Boolean algebra structure.

\(^1\)In fact, it suffices that \( \mathcal{M} \) satisfy the choice principle CP, which is written as follows, in the language of ZF with a new binary relation symbol \( \triangleleft \): “\( \triangleleft \) is a well ordering relation on \( \mathcal{M} \)”.

It is well known that, in every countable model of ZFC, we can define such a binary symbol, so as to get a model of ZF + CP. Thus, ZF + CP is a conservative extension of ZFC.
The language of this model has three binary relation symbols $\in$, $\notin$, and the same function symbols as the model $\mathcal{M}$, with the same interpretation.

The formulas are built as usual, from atomic formulas, with the only logical symbols $\bot$, $\rightarrow$, $\forall$.

We shall use the notations:

- $\forall F$ for $F \rightarrow \bot$; $F_1, \ldots, F_n \rightarrow F$ for $(F_1 \rightarrow (\ldots \rightarrow F_n) \rightarrow F$;
- $\exists x F$ for $\forall \forall x \forall F$; $\exists x(F_1, \ldots, F_n)$ for $\forall \forall x(F_1, \ldots, F_n \rightarrow \bot)$.

**Notation.** We shall often use the notation $\bar{x}$ for a finite sequence $x_1, \ldots, x_n$; for instance, we shall write $F[\bar{x}]$ for $F[x_1, \ldots, x_n]$.

By means of the completeness theorem, we obtain from $\mathcal{N}$ an ordinary model $\mathcal{N}'$, with truth values in $\{0, 1\}$. The set of individuals of $\mathcal{N}'$ generally strictly contains $\mathcal{N}$.

The elements of $\mathcal{N}'$ (resp. $\mathcal{M}$) are called individuals of $\mathcal{N}$ (resp. $\mathcal{M}$). The individuals are generally denoted by $a, b, c, \ldots, a_0, a_1, \ldots$.

In [4] or [5], we define a theory $\text{ZF}_\mathcal{E}$, written in this language. We show that it is a conservative extension of $\text{ZF}$, and that the model $\mathcal{N}$ satisfies the axioms of $\text{ZF}_\mathcal{E}$, which means that each of these axioms is realized by a proof-like term.

Given a term $\xi \in \Lambda$ and a closed formula $F[a_1, \ldots, a_n]$ in the language of $\text{ZF}_\mathcal{E}$, with parameters $a_1, \ldots, a_n$ in $\mathcal{N}$ (or, which is the same, in $\mathcal{M}$), we shall write:

$\xi \models F[a_1, \ldots, a_n]$ in order to say that the term $\xi$ realizes $F[a_1, \ldots, a_n]$.

The truth value of this formula is a subset of $\Pi$, denoted by $\|F[a_1, \ldots, a_n]\|$.

We write $\models F[a_1, \ldots, a_n]$ in order to say that $F[a_1, \ldots, a_n]$ is realized by some proof-like term.

Thus, the model $\mathcal{N}'$ satisfies $\text{ZF}_\mathcal{E}$; therefore, in $\mathcal{N}'$, we can define a model of $\text{ZF}$, denoted $\mathcal{N}'_{\mathcal{E}}$, in which the equality is interpreted by the extensional equivalence denoted by $x = y$ (that is $x \subseteq y \wedge y \subseteq x$).

The general properties of the realizability models are described in [6]; we shall use the definitions and notations of this paper.

In what follows, unless otherwise stated, each formula of $\text{ZF}_\mathcal{E}$ must be interpreted in $\mathcal{N}$ (its truth value is a subset of $\Pi$) or, if one prefers, in $\mathcal{N}'$ (then its truth value is 0 or 1). If the formula must be interpreted in $\mathcal{M}$, (in that case, it does not contains the symbol $\notin$) it will be explicitly stated.

**Function symbols**

**Notations.** The formula $\forall z(z \notin y \rightarrow z \notin x)$ is denoted by $x \subseteq y$ (strong inclusion);
the formula $x \subseteq y \wedge y \subseteq x$ is denoted by $x \equiv y$ (strong extensional equivalence).

We recall that $\subseteq$ and $\equiv$ are the symbols of inclusion and of extensional equivalence of ZF:

$x \subseteq y \equiv \forall z(z \notin y \rightarrow z \notin x)$; $x \equiv y \equiv (x \subseteq y \wedge y \subseteq x)$.

**Function symbols associated with axioms of $\text{ZF}_\mathcal{E}$**

In this section, we define a function symbol for each of the following axioms of $\text{ZF}_\mathcal{E}$:

**Comprehension.**

For each formula $F[y, \bar{z}]$ of $\text{ZF}_\mathcal{E}$, (where $\bar{z}$ is a finite sequence of variables $z_1, \ldots, z_n$) we define,
We show that \( \| \xi \| = \| F[b, \tilde{c}] \| \).

It was shown in [6] (and it is easily checked) that we have:

\[
\| b \not\in \text{Compr}_P(x, \tilde{z}) \| = \| F[y, \tilde{z}] \| \rightarrow y \not\in x \}
\]

Therefore, instead of Compr\(x, \tilde{z}\), we shall use for this function symbol, the more intuitive notation \( \{ y \in x ; F[y, \tilde{z}] \} \), in which \( y \) is a bounded variable.

**Pairing.**

We define the following binary function symbol:

\[
\text{pair}(x, y) = \{ z \in \{ x, y \} \times \Pi ; (z = x) \lor (z = y) \}.
\]

It is easily checked that we have the desired property:

\[
\| \forall x \forall y \forall z \exists \xi \text{ pair}(x, y) \rightarrow z = x \lor z = y \|.
\]

**Remark.** We could also define a symbol pair\(x, y\), with this property, directly in \( \mathcal{M} \), as follows:

\[
\text{pair}(x, y) = \{ (x, \bot \cdot \pi) ; \pi \in \Pi \} \cup \{ (y, \emptyset \cdot \pi) ; \pi \in \Pi \}.
\]

In the sequel, when working in \( \mathcal{M} \), we shall use the (natural) abbreviations:

\[
\{ x, y \} \text{ for pair}(x, y) ; (x, y) \text{ for pair(pair}(x, x), \text{pair}(x, y)).
\]

**Union and power set.**

We define below two unary function symbols \( \bigcup x \) and \( \mathcal{P}(x) \), such that:

\[
\| \forall x \forall z \exists \xi \bigcup x = (\exists y \in x \forall z \in y) \|.
\]

\[
\| \forall x \forall y \bigcup x \lor \exists z \in x \exists \xi \mathcal{P}(x) = \forall y \lor \exists z \in x \exists \xi \lor \forall z \in y \exists \xi \|.
\]

**Theorem 1.** Let \( \mathcal{V}, \mathcal{Q} \) be the unary function symbols defined in \( \mathcal{M} \) as follows:

\[\mathcal{V}(a) = \text{Cl}(a) \times \Pi \text{ and } \mathcal{Q}(a) = \mathcal{P}(\text{Cl}(a) \times \Pi) \times \Pi\]

where Cl\(a\) is the transitive closure of \( a \). Then, we have:

i) \( l \rightarrow \forall x \forall y \forall z \exists \xi \mathcal{V}(x) \rightarrow y \not\in x \}

\[
\text{ii) } l \rightarrow \forall x \forall z \exists \xi \{ y \in x ; F[y, \tilde{z}] \} \lor \mathcal{Q}(x) \text{ for every formula } F[x, \tilde{z}] \text{ of } \mathcal{ZFC}.
\]

i) Let \( a, b, c \) be individuals in \( \mathcal{M} \), \( \xi, \eta \in \Lambda \) and \( \pi \in \Pi \) such that:

\[
\xi \models c \in b, \eta \models c \not\in \mathcal{V}(a) \text{ and } \pi \in \| b \not\in a \| ; \text{ we have therefore } (b, \pi) \in a.
\]

We must show \( \xi \star \eta \cdot \pi \in \perp \).

We show that \( \| c \not\in b \| = \| c \not\in \mathcal{V}(a) \| \): indeed, if \( \rho \in \| c \not\in b \| \), then we have \( (c, \rho) \in b \). But we have \( (b, \pi) \in a \) and thus \( c \in \text{Cl}(a) \) and it follows that \( \| c \not\in \mathcal{V}(a) \| = \Pi \).

Therefore, \( \eta \models c \not\in b \); by hypothesis on \( \xi \), we have \( \xi \star \eta \cdot \pi \in \perp \).

ii) Let \( a, \tilde{c} \) be individuals in \( \mathcal{M} \); we must show \( l \models A \not\in \mathcal{Q}(a) \), where \( A = \{ y \in a ; F[y, \tilde{c}] \} \).

We have \( A = \{ (b, \xi \cdot \pi) ; (b, \pi) \in a, \xi \models F[b, \tilde{c}] \} \) and therefore \( A \in \text{Cl}(a) \times \Pi \). But we have:

\[
\| A \not\in \mathcal{Q}(a) \| = \{ \pi \in \Pi ; (A, \pi) \in \mathcal{Q}(a) \} = \Pi \text{ and therefore } l \models A \not\in \mathcal{Q}(a).
\]

Q.E.D.

We can now define the function symbols \( \bigcup \) and \( \mathcal{P} \) by setting:

\[
\bigcup x = \{ z \in \mathcal{V}(x) ; (\exists y \in x) z \in y \}; \mathcal{P}(x) = \{ y \in \mathcal{Q}(x) ; y \subseteq x \}.
\]

**Collection.**

We shall use in the following, function symbols associated with a strong form of the collection scheme.
In order to define these function symbols, it is convenient to decompose them, which is done in theorems 2, 3 and 4.

Theorem 2. For each formula \( F(x, \vec{z}) \) of \( ZF_\varepsilon \), we have:

\[
\models \forall \vec{z} \left( \exists x F(x, \vec{z}) \rightarrow (\exists x \epsilon \phi_F(\vec{z})) F(x, \vec{z}) \right) ;
\models \forall \vec{z} (\forall x \epsilon \phi_F(\vec{z})) F(x, \vec{z})
\]

where \( \phi_F \) is a function symbol defined in \( \mathcal{M} \).

We show \( \lambda x(\vec{z}) \models \forall x (x \epsilon \Phi_F(\vec{z}) \rightarrow F(x, \vec{z})) \rightarrow \forall x F(x, \vec{z}) \) where the function symbol \( \Phi_F \) is defined as follows:

By means of the collection scheme in \( \mathcal{M} \), we define a function symbol \( \Psi(\vec{z}) \) such that:

\[
\| \forall x F(x, \vec{z}) \| = \bigcup_{x \epsilon \Psi(\vec{z})} \| F(x, \vec{z}) \| \quad \text{and we set} \quad \Phi_F(\vec{z}) = \Psi(\vec{z}) \times \Pi.
\]

Let \( \xi \models \forall x (x \epsilon \Phi_F(\vec{z}) \rightarrow F(x, \vec{z})) \) and \( \pi \in \| \forall x F(x, \vec{z}) \| \).

Then \( \pi \in \| F(x, \vec{z}) \| \) for some \( x \epsilon \Psi(\vec{z}) \), and therefore \( \models \neg x \epsilon \Phi_F(\vec{z}) \) and \( \xi * 1. \pi \epsilon \bot \).

Therefore, by replacing \( F \) with \( \neg F \), we have \( \models \exists x F(x, \vec{z}) \rightarrow (\exists x \epsilon \Phi_F(\vec{z})) F(x, \vec{z}) \).

Thus, we only need to set \( \Phi_F(\vec{z}) = \{ x \epsilon \Phi_{\neg F}(\vec{z}) ; F(x, \vec{z}) \} \).

Q.E.D.

Theorem 3. For every formula \( F(y, \vec{z}) \) of \( ZF_\varepsilon \), we have:

\[
\models \forall \vec{z} \left( \exists x \forall y F(y, \vec{z}) \rightarrow y \epsilon x \right) \rightarrow \forall y (F(y, \vec{z}) \leftrightarrow y \epsilon \gamma_F(\vec{z}))
\]

where \( \gamma_F \) is a function symbol defined in \( \mathcal{M} \).

By theorem 2, we have:

\[
\models \forall \vec{z} \left( \exists x \forall y F(y, \vec{z}) \rightarrow y \epsilon x \right) \rightarrow (\exists x \epsilon \phi(\vec{z})) \forall y (F(y, \vec{z}) \rightarrow y \epsilon x)
\]

where \( \phi \) is a function symbol. Therefore we have, by definition of \( \overline{\bigcup} \phi(\vec{z}) :\)

\[
\models \forall \vec{z} \left( \exists x \forall y F(y, \vec{z}) \rightarrow y \epsilon x \right) \rightarrow \forall y (F(y, \vec{z}) \rightarrow y \epsilon \overline{\bigcup} \phi(\vec{z}))
\]

Now, we only need to set \( \gamma_F(\vec{z}) = \{ y \epsilon \overline{\bigcup} \phi(\vec{z}) ; F(y, \vec{z}) \} \) (comprehension scheme).

Q.E.D.

When the hypothesis \( \exists x \forall y F(y, \vec{z}) \rightarrow y \epsilon x \) is satisfied, we say that the formula \( F(y, \vec{z}) \) defines a set.

For the function symbol \( \gamma_F(\vec{z}) \), we shall use the more intuitive notation \( \{ y ; F(y, \vec{z}) \} \), where \( y \) is a bounded variable.

Theorem 4.

Let \( f(x, \vec{z}) \) be an \( (n + 1) \)-ary function symbol (defined in \( \mathcal{M} \)). Then, we have:

\[
\models \forall a \forall y \forall \vec{z} \left( y \epsilon \phi_f(a, \vec{z}) \leftrightarrow (\exists x \epsilon a)(y = f(x, \vec{z})) \right)
\]

where \( \phi_f \) is an \( (n + 1) \)-ary function symbol.

We define, in \( \mathcal{M} \), the symbol \( \phi_f \) as follows:

Let \( a_0 , y_0 , \vec{z}_0 \) be fixed individuals in \( \mathcal{M} \); we set \( \phi_f(a_0 , \vec{z}_0) = \{(f(x, \vec{z}_0), y) ; (x, y) \epsilon a_0 \} \).

Then, we have immediately \( \| y \epsilon \phi_f(a_0 , \vec{z}_0) \| = \| \forall x (y_0 = f(x, \vec{z}_0) \rightarrow x \epsilon a_0) \| \).

Therefore:

\[
\models \forall x (y_0 = f(x, \vec{z}_0) \rightarrow x \epsilon a_0) \Rightarrow y_0 \epsilon \phi_f(a_0 , \vec{z}_0)
\]

which gives the desired result.

Q.E.D.

Remark. The connective \( \rightarrow \) is defined in [5, 6].

For the function symbol \( \phi_f(a, \vec{z}) \), we shall use the more intuitive notation \( \{ f(x, \vec{z}) ; x \epsilon a \} \), where \( x \) is a bounded variable. We call it \( \text{image of } a \text{ by the function } f(x) \).
Symbols for characteristic functions

Let \( R(x_1, \ldots, x_n) \) be a \( n \)-ary relation defined in \( \mathcal{M} \). Its characteristic function, with values in \([0, 1]\), will be denoted by \( \langle R(x_1, \ldots, x_n) \rangle \). Therefore, we have: \( \mathcal{M} \models \forall \vec{x} (R(\vec{x}) \leftrightarrow \langle R(\vec{x}) \rangle = 1) \). Therefore, in the realizability model \( \mathcal{N} \), the function symbol \( \langle R(\vec{x}) \rangle \) takes its values in \( \mathbb{2} \).

The theorem 8 below, shows that, if a binary relation \( y \prec x \) is well founded in \( \mathcal{M} \), then the relation \( \langle y \prec x \rangle = 1 \) is well founded in \( \mathcal{N} \).

Miscellaneous symbols

In the following, we shall use some function symbols, the definition and properties of which are given in [6]. We simply recall their definition below.

- The unary function symbol \( \downarrow \), defined in \( \mathcal{M} \) by \( \downarrow x = x \times \Pi \).
  For any individual \( E \) of \( \mathcal{M} \), the restricted quantifier \( \forall x \downarrow E \) is defined in [5] or [6] by: \( \| \forall x \downarrow E F[x] \| = \bigcup_{x \in E} \| F[x] \| \) and we have \( \| - \forall x \downarrow E F[x] \| \leftrightarrow \forall x (x \in E \rightarrow F[x]) \).
  In the realizability model \( \mathcal{N} \), the formula \( x \in E \) may be intuitively understood as “\( x \) is of type \( E \)”. For instance, \( \mathbb{2} \) may be considered as the type of booleans and \( \mathbb{N} \) as the type of integers.

- The function symbols \( \land, \lor, \neg \), with domains \([0, 1] \times [0, 1] \) and \([0, 1] \), and values in \([0, 1]\), are defined in \( \mathcal{M} \) by means of the usual truth tables.
  These functions define, in \( \mathcal{N} \), a structure of Boolean algebra on \( \mathbb{2} \).
  We call it the characteristic Boolean algebra of the realizability model \( \mathcal{N} \).

- A binary function symbol with domain \([0, 1] \times \mathcal{M} \), denoted by \( (\alpha, x) \mapsto \alpha x \), by setting: \( 0x = \varnothing ; 1x = x \).
  In the model \( \mathcal{N} \), the domain of this function is \( \mathbb{2} \times \mathcal{N} \).

- A binary function symbol \( \sqcup \) with domain \( \mathcal{M} \times \mathcal{M} \), by setting \( x \sqcup y = x \sqcup y \).
  \textbf{Remark.} The extension of this function to the model \( \mathcal{N} \) is not the union \( \sqcup \), which explains the use of another symbol.

\textbf{Lemma 5 (Linearity).}

Let \( f \) be a binary function symbol, defined in \( \mathcal{M} \). Then, we have:

i) \( l \models \forall \alpha \mathbb{2} \forall x \forall y (\alpha f(x, y) = \alpha f(\alpha x, y)) \).

ii) Moreover, if \( f(\varnothing, \varnothing) = \varnothing \), then:

\( l \models \forall \alpha \mathbb{2} \forall \alpha' \mathbb{2} \forall x \forall y \forall x' \forall y' (\alpha \land \alpha' = 0 \mapsto f(\alpha x \sqcup \alpha' x', \alpha y \sqcup \alpha' y') = \alpha f(x, y) \sqcup \alpha' f(x', y')) \).

\textbf{Remark.} The connective \( \mapsto \) is defined in [5, 6].

It suffices to check:

for (i) the two cases \( \alpha = 0, 1 \);

for (ii) the three cases \( (\alpha, \alpha') = (0, 0), (0, 1), (1, 0) \);

which is is trivial.

Q.E.D.
Well founded relations

In this section, we study properties of well founded relations in $\mathcal{N}$. All the results obtained here are, of course, trivial in ZF. The difficulties come from the fact that the relation $\epsilon$ of strong membership, does not satisfy extensionality.

Given a binary relation $\prec$, an individual $a$ is said minimal for $\prec$, if we have $\forall x \prec (x < a)$.

The binary relation $\prec$ is called well founded if we have:

$$\forall X \left( \forall x (\forall y (y < x \rightarrow y \notin X) \rightarrow x \notin X) \rightarrow \forall x (x \notin X) \right).$$

The intuitive meaning is that each non-empty individual $X$ has an $\epsilon$-element minimal for $\prec$.

**Theorem 6.**

If the relation $x < y$ is well founded then, for every formula $F[x, \bar{z}]$ of ZF, we have:

$$\forall \bar{z} \left( \forall x (\forall y (y < x \rightarrow F[y, \bar{z}]) \rightarrow F[x, \bar{z}] \rightarrow \forall x F[x, \bar{z}] \right).$$

Proof by contradiction; we consider, in $\mathcal{N}$, an individual $a$ and a formula $G[x]$ such that:

1. $G[a] ; \forall x \{ G[x] \rightarrow \exists y (G[y], y < x) \}.$

We apply the axiom scheme of infinity of ZF:

2. $\exists c \{ a \in c, (\forall x \in c) \{ \exists y F(x, y) \rightarrow (\exists y \epsilon c)F(x, y) \} \}$

by setting $F(x, y) \equiv G[x] \land G[y] \land y < x$. Let $b = \{ x \epsilon c ; G(x) \}$ ; by (1) and (2), we get $a \epsilon b$.

We obtain a contradiction with the hypothesis, by showing $(\forall x \epsilon b)(\exists y \epsilon b)(y < x)$.

Therefore, we suppose $x \epsilon c$ and $G[x]$; by (2), we have:

$\exists y(G[x], G[y], y < x) \rightarrow (\exists y \epsilon c)\{ G[x], G[y], y < x \}.$

By $G[x]$ and (1), we have $\exists y G[x], G[y], y < x$.

Therefore, we have $(\exists y \epsilon c)(G[y], y < x)$, hence the result.

Q.E.D.

Therefore, in order to show $\forall x F[x]$, it suffices to show $\forall x (\forall y (y < x \rightarrow F[y]) \rightarrow F[x])$.

Then, we say that we have shown $\forall x F[x]$ by induction on $x$, following the well founded relation $\prec$.

**Theorem 7.** The binary relation $x \epsilon y$ is well founded.

We must show $\forall x (\forall y (y \epsilon x \rightarrow y \notin X) \rightarrow x \notin X) \rightarrow \forall x (x \notin X)$.

We apply theorem 6 to the well founded relation $x \epsilon y$ and the formula $F[x] \equiv x \notin X$.

This gives: $\forall x (\forall y (y \epsilon x \rightarrow y \notin X) \rightarrow x \notin X) \rightarrow \forall x (x \notin X)$.

Now, we have immediately $|\vdash x \notin X \rightarrow x \notin X$. Thus, it remains to show:

$|\vdash \forall x (\forall y (y \epsilon x \rightarrow y \notin X) \rightarrow x \notin X) \rightarrow \forall x (\forall y (y \epsilon x \rightarrow y \notin X) \rightarrow x \notin X)$.

But we have $x \in X \equiv \forall x' (x' \approx x \rightarrow x' \notin X)$. Therefore, we need to show:

$|\vdash \forall x (\forall y (y \epsilon x \rightarrow y \notin X) \rightarrow x \notin X), \forall y (y \epsilon x \rightarrow y \notin X), x' \approx x \rightarrow x' \notin X$; or else:

$|\vdash y (y \epsilon x \rightarrow y \notin X), x' \approx x \rightarrow \forall y (y \epsilon x' \rightarrow y \notin X)$.

Now, from $x' \approx x$, we deduce $y \in x$. Thus, there is some $y' \approx y$ such that $y' \epsilon x$.

Then, from $\forall y (y \epsilon x \rightarrow y \notin X)$, we deduce $y' \notin X$, and therefore $y \notin X$.

Q.E.D.

For instance, in the following, we shall use the fact that, if there is an ordinal $\rho$ such that $F[\rho]$, then there exists a least such ordinal, for any formula $F[\rho]$ written in the language of ZF. This results from theorem 7.
Preservation of well-foundedness

**Theorem 8.** Let $<$ be a well founded binary relation, defined in the ground model $\mathcal{M}$. Then, the relation $\langle y < x \rangle = 1$ is well founded in $\mathcal{N}$. In fact, we have:
\[
\mathcal{N} \models \forall X (\forall y (\forall x(y < x) = 1 \rightarrow y \not\in X) \rightarrow \forall x (x \not\in X))
\]
where $Y = (\lambda x \lambda f(x)f)(x)xf$ (Turing fixpoint combinator).

Let $\xi \in \Lambda$ be such that $\xi \models \forall y(\forall x(y < x) = 1 \rightarrow y \not\in X_0) \rightarrow x \not\in X_0$, $X_0$ being any individual in $\mathcal{M}$. We set $F[x] \equiv (\forall \pi \in \|x \not\in X_0\|) (Y \star \xi \star \pi \in \perp)$, and we have to show $\forall x F[x]$.

Since $<$ is a well founded relation, it suffices to show $\forall x (\forall y(y < x \rightarrow F[y]) \rightarrow F[x])$, or else $\neg F[x_0] \rightarrow (\exists y < x_0) \neg F[y]$, for any individual $x_0$.

By the hypothesis $\neg F[x_0]$, there exists $\pi_0 \in \|x_0 \not\in X_0\|$ such that $Y \star \xi \star \pi_0 \not\in \perp$ and therefore, we have $\xi \star Y \xi \star \pi_0 \not\in \perp$.

By hypothesis on $\xi$, we deduce $Y \xi \not\models \forall y(\forall x(y < x_0) = 1 \rightarrow y \not\in X_0)$.

Thus, there exists $y_0 < x_0$ such that $Y \xi \not\models y_0 \not\in X_0$.

Therefore, we have $(\exists \pi \in \|y_0 \not\in X_0\|)(Y \star \xi \star \pi \not\in \perp)$, that is $\neg F[y_0]$.

Q.E.D.

**Definition of a rank function**

**Definition.** A function with domain $D$ is an individual $\phi$ such that:

\[
(\forall z \in \phi)(\exists x \in D) \exists y(x = (x, y)) ; (\forall x \in D) \exists y((x, y) \in \phi) ;
\]

\[
\forall x \forall y \forall y'( (x, y) \in \phi, (x, y') \in \phi \rightarrow y = y').
\]

Let $\phi$ be a function with domain $D$ and $F[y, \vec{z}]$ a formula of $\text{ZF}_c$. Then, the formula:

\[
\exists y((x, y) \in \phi, F[y, \vec{z}])
\]

is denoted by $F[\phi(x), \vec{z}]$.

**Remark.** Beware, despite the same notation $\phi(x)$, it is not a function symbol.

By means of theorem 3, we define the binary function symbol $\text{Im}$ by setting:

\[
\text{Im}(\phi, D) = \{y ; (\exists x \in D) (x, y) \in \phi\}.
\]

When $\phi$ is a function with domain $D$, we shall use, for $\text{Im}(\phi, D)$, the more intuitive notation $\{\phi(x) ; x \in D\}$, which we call image of the function $\phi$.

Let $D' \subseteq D$, that is $\forall x (x \in D \rightarrow x \in D')$; a restriction of $\phi$ to $D'$ is, by definition, a function $\phi'$ with domain $D'$ such that $\phi \subseteq \phi'$.

For instance, $\{z \in \phi ; (\exists x \in D') \exists y(x = (x, y))\}$ is a restriction of $\phi$ to $D'$.

If $\phi_0', \phi_1'$ are both restrictions of $\phi$ to $D'$, then $\phi_0' \equiv \phi_1'$.

**Definition.**

A binary relation $<$ is called ranked, if we have $\forall x \exists y \forall z (z < x \rightarrow z \in y)$, in other words: the minorants of any individual form a set.

By theorem 3, if the relation $<$ is ranked and defined by a formula $P[x, y, \vec{u}]$ of $\text{ZF}_c$ with parameters $\vec{u}$ in $\mathcal{N}$, we have:

\[
\mathcal{N} \models \forall x \forall y (x < y \leftrightarrow x \in f(y, \vec{u})) ,
\]

for some symbol of function $f$, defined in $\mathcal{M}$.

In what follows, we suppose that $<$ is a ranked transitive binary relation.

A function $\phi$ with domain $\{x ; x < a\}$ will be called $a$-inductive for $<$, if we have:

\[
\phi(x) = \{\phi(y) ; y < x\} \text{ for every } x < a.
\]

In other words:

\[
(\forall x < a)(\forall y < x) \phi(y) \in \phi(x) ; (\forall x < a)(\forall y \in \phi(x))(\exists y < x) z \in \phi(y).
\]

If $\phi$ is $a$-inductive for $<$, we set $O(\phi, a) = \{\phi(x) ; x < a\}$ (image of $\phi$).
Lemma 9. Let \( \phi, \phi' \) be two functions, \( a \)-inductive for \( < \). Then:

i) \( \phi(x) \simeq \phi'(x) \) for every \( x < a \).

ii) \( O(\phi, a) \simeq O(\phi', a) \).

iii) \( (\forall x < a) On(\phi(x)) \); \( O(\phi, a) \) is an ordinal, called ordinal of \( \phi \).

i) Proof by induction on \( \phi(x) \), following \( \in \) : if \( u \in \phi(x) \), then \( u \simeq \phi(y) \) with \( y < x \). Since \( \phi(y) \in \phi(x) \), we have \( \phi(y) = \phi'(y) \) by the induction hypothesis; therefore \( \phi(y) \in \phi'(x) \) and \( \phi(x) \in \phi'(x) \).

Conversely, if \( u \in \phi'(x) \), then \( u \simeq \phi'(y) \) with \( y < x \). Thus, we have \( \phi(y) \in \phi(x) \), and therefore \( \phi(y) \simeq \phi'(y) \) by the induction hypothesis; therefore \( u \in \phi(x) \) and \( \phi'(x) \subset \phi(x) \).

ii) Immediate, by (i).

iii) We show \( On(\phi(x)) \) by induction on \( \phi(x) \), for the well founded relation \( \in \) :

If \( u \in \phi(x) \), we have \( u \simeq \phi(y) \) with \( y < x \); therefore, we have \( On(u) \) by the induction hypothesis. If \( u \in u \), then \( u \in \phi(y) \), therefore \( v = \phi(z) \) with \( z < y \); therefore \( v \in \phi(x) \).

It follows that \( \phi(x) \) is a transitive set of ordinals, thus an ordinal.

Then, \( O(\phi, a) \) is also a transitive set of ordinals, and therefore an ordinal.

Q.E.D.

Lemma 10. If \( \phi \) is \( a \)-inductive for \( < \), and if \( b < a \), then every restriction \( \psi \) of \( \phi \) to the domain \( \{ x ; x < b \} \) is a \( b \)-inductive function for \( < \).

Indeed, we have, \( \psi(x) = \phi(x) \simeq \{ \phi(y) ; y < x \} \simeq \{ \psi(y) ; y < x \} \).

Q.E.D.

By means of theorem 2, we define a unary function symbol \( \Phi \), such that:

\[ \forall x(\forall f \in \Phi(x)) (f \text{ is a } x \text{-inductive function}) ; \]

\[ \forall x \forall f \left( f \text{ is a } x \text{-inductive function } \rightarrow \exists f \in \Phi(x) \right) . \]

In other words, \( \Phi(x) \) is a set of \( x \)-inductive functions, which is non void if there exists at least one such function.

Finally, we define the unary function symbol \( R_k \), using theorem 4, by setting:

\[ R_k(x) = \bigcup \{ O(f, x) ; f \in \Phi(x) \} \]

(the symbol \( \bigcup \) is defined after theorem 1).

Therefore, \( R_k(x) \) is the union of the ordinals of the \( x \)-inductive functions in the set \( \Phi(x) \).

Since all these ordinals are extensionally equivalent, by lemma 9(ii), their union \( R_k(x) \) is also an equivalent ordinal.

Remarks.

If there exists no \( x \)-inductive function, then \( R_k(x) \) is void.

The function symbols \( O, \Phi, R_k \) have additional arguments, which are the parameters \( \bar{u} \) of the formula \( P[x, y, \bar{u}] \) which defines the relation \( y < x \).

We suppose now that \( < \) is a ranked transitive relation, which is well founded. It is therefore a strict ordering.

Lemma 11. Every restriction of \( R_k \) to the domain \( \{ x ; x < a \} \) is an \( a \)-inductive function for \( < \).

Proof by induction on \( a \), following \( < \).

Let \( f \) be a restriction of \( R_k \) to the domain \( \{ x ; x < a \} \) and let \( x < a \). We must show that
\[ f(x) = \{ f(y) ; y < x \}, \text{ in other words, that we have:} \]
\[ \text{Rk}(x) = \{ \text{Rk}(y) ; y < x \}. \]

Let \( \psi \) be any restriction of \( \text{Rk} \) to the domain \( \{ y ; y < x \} \). By the induction hypothesis, \( \psi \) is a \( x \)-inductive function for \( < \).

We now show that \( \text{Rk}(x) = \{ \text{Rk}(y) ; y < x \} \):

i) If \( u \in \text{Rk}(x) \) then \( u \in O(\phi, x) \) for some function \( \phi \) which is \( x \)-inductive for \( < \), provided that there exists such a function. Now, there exists effectively one, otherwise \( \text{Rk}(x) \) would be void.

Therefore, by definition of \( O(\phi, x) \), we have \( u = \phi(y) \) with \( y < x \). But \( \text{Rk}(y) = \phi(y) \), since \( \phi, \psi \) are both \( x \)-inductive functions for \( < \), and \( \psi(y) = \text{Rk}(y) \) (lemma 9(i)).

Therefore, we have \( u \equiv \text{Rk}(y) \), with \( y < x \).

ii) Conversely, if \( y < x \), then \( \text{Rk}(y) = \psi(y) \). Let \( \phi \in \Phi(x) \); then \( \phi, \psi \) are \( x \)-inductive for \( < \); therefore \( \phi(y) \equiv \psi(y) \) (lemma 9(i)).

Now \( \phi(y) \in O(\phi, x) \), and therefore \( \phi(y) \in \text{Rk}(x) \) by definition of \( \text{Rk}(x) \).

It follows that \( \text{Rk}(y) = \psi(y) \in \text{Rk}(x) \).

Q.E.D.

**Theorem 12.** We have \( \text{Rk}(x) = \{ \text{Rk}(y) ; y < x \} \) for every \( x \).

Proof by induction on \( x \), following \( < \); let \( \psi \) be any restriction of \( \text{Rk} \) to the domain \( \{ y ; y < x \} \).

By lemma 11, \( \psi \) is a \( x \)-inductive function for \( < \).

Then, we finish the proof, by repeating paragraphs (i) and (ii) of the proof of lemma 11.

Q.E.D.

\( \text{Rk} \) is called the rank function of the ranked, well founded and transitive relation \( < \).

\( \text{Rk}(x) \) is, for every \( x \), a representative of the ordinal of any \( x \)-inductive function for \( < \).

The values of the rank function \( \text{Rk} \) form an initial segment of \( \text{On} \), which we shall call the image of \( \text{Rk} \).

It is therefore, either an ordinal, or the whole of \( \text{On} \).

**Proposition 13.** Let \( <\_0, <\_1 \) be two ranked transitive well founded relations, and \( f \) a function such that \( \forall x \forall y (x <\_0 y \rightarrow f(x) <\_1 f(y)) \).

If \( \text{Rk}_0, \text{Rk}_1 \) are their rank functions, then we have \( \forall x (\text{Rk}_0(x) \leq \text{Rk}_1(f(x))) \), and the image of \( \text{Rk}_0 \) is an initial segment of the image of \( \text{Rk}_1 \).

We show immediately \( \forall x (\text{Rk}_0(x) \leq \text{Rk}_1(f(x))) \) by induction following \( <\_0 \). Hence the result, since the image of a rank function is an initial segment of \( \text{On} \).

Q.E.D.

**An ultrafilter on \( \mathcal{J}_2 \)**

In all of the following, we write \( y < x \) for \( y \in \text{Cl}(x) \) in \( \mathcal{M} \), where \( \text{Cl}(x) \) denotes the transitive closure of \( x \). It is a strict well founded ordering (many other such orderings would do the job, for instance the relation \( \text{rank}(y) < \text{rank}(x) \)).

The binary function symbol \( (y < x) \) is therefore defined in \( \mathcal{N} \), with values in \( \mathcal{J}_2 \).

By theorem 8, the binary relation \( \langle y < x \rangle = 1 \) is well founded in \( \mathcal{N} \).

**Theorem 14.** \( \models \text{There exists an ultrafilter } \mathcal{D} \text{ on } \mathcal{J}_2 \), which is defined as follows:

\[ \mathcal{D} = \{ \alpha \in \mathcal{J}_2 ; \text{the relation } \langle y < x \rangle \geq \alpha \text{ is well founded} \}. \]
Remark. By lemma 5, the formula \( \langle y < x \rangle \) may be written \( \alpha y < \alpha x = \alpha \).

The formula \( \alpha \in \emptyset \), which we shall also write \( \emptyset \langle \alpha \rangle \), is therefore:

\[
\emptyset \langle \alpha \rangle \equiv \forall X (\forall x (\forall y ((y < x) \geq \alpha \rightarrow \langle y \in X \rangle \rightarrow \forall x (\forall y (x \in X)))
\]

Remark. By lemma 5, the formula \( x \in X \). We have immediately:

\[
\emptyset \langle \alpha \rangle \equiv \forall X (\forall x (\forall y ((y < x) \geq \alpha \rightarrow \langle y \in X \rangle \rightarrow \forall x (\forall y (x \in X))).
\]

Therefore, in order to prove theorem 14, it suffices to show:

\[
\emptyset \langle \alpha \rangle \rightarrow \emptyset \langle \alpha \rangle \equiv \forall X (\forall y ((y < x) \geq \alpha \rightarrow \langle y \in X \rangle \rightarrow \forall x (\forall y (x \in X))).
\]

We have immediately:

\[
\emptyset \langle \alpha \rangle \rightarrow \emptyset \langle \alpha \rangle \equiv \forall X (\forall y ((y < x) \geq \alpha \rightarrow \langle y \in X \rangle \rightarrow \forall x (\forall y (x \in X))).
\]

Therefore, in order to prove theorem 14, it suffices to show:

\[
\emptyset \langle \alpha \rangle \rightarrow \emptyset \langle \alpha \rangle \equiv \forall X (\forall y ((y < x) \geq \alpha \rightarrow \langle y \in X \rangle \rightarrow \forall x (\forall y (x \in X))).
\]

Remark. For \( \alpha \in \emptyset \), we shall write \( x <_\alpha y \) for \( \langle x < y \rangle \geq \alpha \).

**Theorem 15.**

i) \( \emptyset \langle \alpha \rangle \rightarrow \emptyset \langle \alpha \rangle \equiv \forall X (\forall y ((y < x) \geq \alpha \rightarrow \langle y \in X \rangle \rightarrow \forall x (\forall y (x \in X))).
\]

ii) \( \emptyset \langle \alpha \rangle \rightarrow \emptyset \langle \alpha \rangle \equiv \forall X (\forall y ((y < x) \geq \alpha \rightarrow \langle y \in X \rangle \rightarrow \forall x (\forall y (x \in X))).
\]

i) Let \( \alpha, \beta \in \emptyset \) be such that \( \alpha \land \beta = 0, \neg \emptyset \langle \alpha \rangle, \neg \emptyset \langle \beta \rangle \). We have to show \( \neg \emptyset \langle \alpha \land \beta \rangle \).

By hypothesis on \( \alpha \) and \( \beta \), there exists individuals \( a_0, A \) (resp. \( b_0, B \)) such that \( a_0 \in \emptyset (\alpha \land \beta) \). Let \( c \in A, c = \alpha a_0 \lor \beta b_0 \). By hypothesis on \( \alpha, \beta \), there exists \( a' \in A \) and \( b' \in B \) such that \( a' <_\alpha a, b' <_\beta b \). If we set \( c' = a' \lor b' \), we have \( c' \in C \), as needed. We also have:

\[
\langle c' = a' \rangle \geq \alpha, \langle a' < a \rangle \geq \alpha, \langle c = a \rangle \geq \alpha \text{; it follows that } \langle c' < c \rangle \geq \alpha.
\]

In the same way, we have \( \langle c' < c \rangle \geq \beta \) and therefore, finally, \( \langle c' < c \rangle \geq \alpha \land \beta \).

ii) We set \( \beta' = \beta \land (\neg \alpha) \); we have \( \alpha \land \beta' = 0 \) and \( \alpha \lor \beta' = \alpha \lor \beta \). Therefore, we have:

\[
\emptyset \langle \alpha \lor \beta \rangle \rightarrow \emptyset \langle \alpha \rangle \lor \emptyset \langle \beta \rangle \text{. Now, we have } \beta' \leq \beta \text{ and therefore } \emptyset \langle \beta' \rangle \rightarrow \emptyset \langle \beta \rangle.
\]

Q.E.D.

**Lemma 16.**

i) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

ii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iv) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

Let \( a, b \) be two individuals.

i) Let \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

ii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iv) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

Let \( a, b \) be two individuals.

i) Let \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

ii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iv) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

Let \( a, b \) be two individuals.

i) Let \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

ii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iii) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

iv) \( \emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \).

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]

\[
\emptyset \langle \alpha < b \rangle \neq 1 \rightarrow \emptyset \langle \alpha < y \rangle \text{.}
\]
iv) Indeed, if \( a, b \) are individuals of \( \mathcal{M} \), we have trivially:
\[
\| \langle a < b \rangle \neq 1 \| = \| a \not\in \text{Cl}(b) \|.
\]
Q.E.D.

**Lemma 17.** The well founded relation \( \langle x < y \rangle = 1 \) is ranked, and its rank function \( R \) has for image the whole of \( \text{On} \).

Lemma 16(iv) shows that this relation is ranked.

Let \( \rho \) be an ordinal and \( r \) an individual \( \equiv \rho \). We show, by induction on \( \rho \), that \( R(r) \geq \rho \).

Indeed, for every \( \rho' \in \rho \), there exists \( r' \in r \) such that \( r' \equiv \rho' \). We have \( R(r') \geq \rho' \) by induction hypothesis, and \( \langle r' < r \rangle = 1 \) from lemma 16(i). Therefore, we have \( \rho' \in R(r) \) by definition of \( R \), and finally \( R(r) \geq \rho \). This shows that the image of \( R \) is not bounded in \( \text{On} \). Since it is an initial segment, it is the whole of \( \text{On} \).

Q.E.D.

**Theorem 18.** Let \( F(x, y) \) be a formula of ZF\(_e\), with parameters. Then, we have:
\[
I \models \forall x \forall y \left( \forall \omega \exists ! F(x, f(x, \omega)) \rightarrow F(x, y) \right)
\]
for some function symbol \( f \), defined dans \( \mathcal{M} \), with domain \( \mathcal{M} \times \Pi \).

Since the ground model \( \mathcal{M} \) satisfies \( V = L \) (or only the choice principle), we can define, in \( \mathcal{M} \), a function symbol \( f \) such that:
\[
\forall x \forall y (\forall \omega \in \Pi \left( \omega \in \| F(x, y) \| \rightarrow \omega \in \| F(x, f(x, \omega)) \| \right)).
\]

Let \( a, b \) be individuals, \( \xi \| \exists \omega \in \Pi F(a, f(a, \omega)) \) and \( \pi \in \| F(a, b) \| \).

Thus, we have \( \pi \in \| F(a, f(a, \pi)) \| \), and therefore \( \xi \parallel \pi \in \bot \).

Q.E.D.

**Definitions.** Let \( a \) be any individual of \( \mathcal{N} \) and \( \kappa \) an ordinal (therefore, \( \kappa \) is not an individual of \( \mathcal{N} \), but an equivalence class for \( \equiv \)).

A function or application from \( \kappa \) into \( a \) is, by definition, a binary relation \( R(\rho, x) \) such that:
\[
\forall x \forall x' (\forall \rho, \rho' \in \kappa \left( R(\rho, x), R(\rho', x'), \rho \equiv \rho' \rightarrow x = x' \right))\} (\forall \rho \in \kappa) (\exists x \in a) R(\rho, x).
\]

It is an injection if we have \( \forall x (\forall \rho, \rho' \in \kappa \left( R(\rho, x), R(\rho', x) \rightarrow \rho \equiv \rho' \right)) \).

A surjection from \( a \) onto \( \kappa \) is a function \( f \) of domain \( a \) such that:
\[
(\forall \rho \in \kappa) (\exists x \in a) f(x) \equiv \rho.
\]

**Theorem 19.** For any individual \( a \), there exists an ordinal \( \kappa \), such that there is no surjection from \( a \) onto \( \kappa \).

Let \( f \) be a surjection from \( a \) onto an ordinal \( \rho \). We define a strict ordering relation \( \prec_f \) by setting \( x <_f y \equiv x \in a \land y \in a \land f(x) < f(y) \). It is clear that this relation is well founded, that \( f \) is an \( a \)-inductive function, and that \( O(f, a) \equiv \rho \).

We may consider this relation as a subset of \( a \times a \).

By means of the axioms of union, power set and collection given above (theorems 1 à 4), we define an ordinal \( \kappa_0 \), which is the union of the \( O(f, a) \) for all the functions \( f \) which are \( a \)-inductive for some well founded strict ordering relation on \( a \).

In fact, we consider the set:
\[
\mathcal{B}(a) = \{ X \in \mathcal{P}(a \times a) ; X \text{ is a well founded strict ordering relation on } a \}.
\]

Then, we set \( \kappa_0 = \bigcup \{ O(f, a) ; X \in \mathcal{B}(a), f \in \Phi(X, a) \} \).

In this definition, we use the function symbol \( \Phi \), defined after lemma 10, which associates
with each well founded strict ordering relation $X$ on $a$, a non void set of $a$-inductive functions for this relation.

Then, there exists no surjection from $a$ onto $\kappa_0 + 1$.

Q.E.D.

We denote by $\Delta$ the first ordinal of $\mathcal{N}$ such that there is no surjection from $\mathfrak{M}$ onto $\Delta$ : for every function $\phi$, there exists $\delta \in \Delta$ such that $\forall x^{\mathfrak{M}}(\phi(x) \neq \delta)$.

For each $\alpha \varepsilon \mathfrak{M}$, we denote by $\mathcal{N}_\alpha$ the class defined by the formula $x = \alpha x$.

**Lemma 20.** Let $\alpha_0, \alpha_1 \varepsilon \mathfrak{M}$, $\alpha_0 = 0$ and $R_0$ (resp. $R_1$) be a functional relation of domain $\mathcal{N}_{\alpha_0}$ (resp. $\mathcal{N}_{\alpha_1}$) with values in $\mathbb{On}$. Then, either $R_0$, or $R_1$, is not surjective onto $\Delta$.

Proof by contradiction : we suppose that $R_0$ and $R_1$ are both surjective onto $\Delta$.

We apply theorem 18 to the formula $F(x_0, x_1) \equiv \neg (R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 x_1))$, and we get :

$$\forall x_0 (\exists x_1(R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 x_1)) \rightarrow \exists x^{\mathfrak{M}}(R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 f(x_0, \omega))))$$

where $f$ is a suitable function symbol (therefore defined in $\mathcal{M}$).

Replacing $x_0$ with $\alpha_0 x_0$, we obtain :

$$\forall x_0 (\exists x_1(R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 x_1)) \rightarrow \exists x^{\mathfrak{M}}(R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 f(\alpha_0 x_0, \omega))))$$

But, by lemma 5(i), we have $\alpha_1 f(\alpha_0 x, \phi) = \alpha_1 f(\alpha_0 x, \omega) = \alpha_1 f(\phi, \omega)$. It follows that :

$$\forall x_0 (\exists x_1(R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 x_1)) \rightarrow \exists x^{\mathfrak{M}}(R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 f(\phi, \omega))))$$

By hypothesis, we have $(\forall \rho \in \Delta) \exists x_0 \exists x_1 (\rho \equiv R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 x_1))$. It follows that :

$(\forall \rho \in \Delta) \exists x_0 \exists x_1 (\rho \equiv R_0(\alpha_0 x_0) \equiv R_1(\alpha_1 f(\phi, \omega)))$ ; therefore, we have :

$(\forall \rho \in \Delta) \exists x_0 \exists x_1 (\rho \equiv R_1(\alpha_1 f(\phi, \omega)))$.

Therefore, the function $\omega \mapsto R_1(\alpha_1 f(\phi, \omega)))$ is a surjection from $\mathfrak{M}$ onto $\Delta$. But this is a contradiction with the definition of $\Delta$.

**Remark.** We should write $f(\alpha_0, \alpha_1, x_0, \omega)$ instead of $f(x_0, \omega)$, since the function symbol $f$ depends on the four variables $\alpha_0, \alpha_1, x_0, \omega$. In fact, it depends also on the parameters which appear in $R_0, R_1$. The proof does not change.

Q.E.D.

**Corollary 21.** Let $\alpha_0, \alpha_1 \varepsilon \mathfrak{M}$, $\alpha_0 \wedge \alpha_1 = 0$, and $<_{0}, <_{1}$ be two well founded ranked strict ordering relations with respective domains $\mathcal{N}_{\alpha_0}, \mathcal{N}_{\alpha_1}$. Let $Rk_0, Rk_1$ be their rank functions. Then, either the image of $Rk_0$, or that of $Rk_1$ is an ordinal $< \Delta$.

In order to be able to define the rank functions $Rk_0, Rk_1$, we consider the relations $\prec_{0}, \prec_{1}$, with domain the whole of $\mathcal{N}$, defined by $x \prec_{i} y \equiv (x = \alpha_i x) \wedge (y = \alpha_i y) \wedge (x <_{i} y)$ for $i = 0, 1$.

These strict ordering relations are well founded and ranked.

Their rank functions $Rk_0', Rk_1'$ take the value $0$ outside $\mathcal{N}_{\alpha_0}, \mathcal{N}_{\alpha_1}$ respectively : indeed, all the individuals outside $\mathcal{N}_{\alpha_i}$ are minimal for $\prec_{i}$.

By lemma 20, one of them, $Rk_0'$ for instance, is not surjective onto $\Delta$.

Since the image of any rank function is an initial segment of $\mathbb{On}$, the image of $Rk_0$ is an ordinal $< \Delta$.

Q.E.D.

**Theorem 22.**

i) $\parallel \forall \alpha_0 \forall \alpha_1 (\alpha_0 \wedge \alpha_1 = 0 \rightarrow (\mathcal{D}[\alpha_0], \mathcal{D}[\alpha_1] \twoheadrightarrow \perp))$.

ii) $\parallel \forall \alpha_0 \forall \alpha_1 (\mathcal{D}[\alpha_0], \mathcal{D}[\alpha_1] \twoheadrightarrow \mathcal{D}[\alpha_0 \wedge \alpha_1])$. 
i) In $\mathcal{N}$, let $\alpha_0, \alpha_1 \in \mathbb{2}$ be such that $\alpha_0 \wedge \alpha_1 = 0$ and the relations $\langle x < y \rangle \geq \alpha_0$, $\langle x < y \rangle \geq \alpha_1$ be well founded. Therefore, we have $\alpha_0, \alpha_1 \neq 0, 1$.

Therefore, the relations $x < _i y \equiv (x = \alpha_i x) \wedge (y = \alpha_i y) \wedge (\langle x < y \rangle = \alpha_i)$ for $i = 0, 1$, are well founded strict orderings.

From lemma 16(iii), it follows that these relations are ranked.

Now, by lemma 5, we have $\models \forall x \forall y \exists \alpha \in \mathbb{2} (\langle x < y \rangle = 1 \rightarrow \langle ax < ay \rangle = \alpha)$.

But, by lemma 17, the rank function of the well founded relation $\langle x < y \rangle = 1$ has for image the whole of $\text{On}$. Therefore, by proposition 13, the same is true for the rank functions of the well founded strict order relations $x < _0 y$ and $x < _1 y$.

But this contradicts corollary 21.

ii) We have $\alpha_0 \leq (\alpha_0 \wedge \alpha_1) \cup (\neg \alpha_1)$. Therefore, by $\mathcal{D}[\alpha_0]$ and theorem 15, we have $\mathcal{D}[\alpha_0 \wedge \alpha_1]$ or $\mathcal{D}[\neg \alpha_1]$. But $\mathcal{D}[\neg \alpha_1]$ is impossible, by $\mathcal{D}[\alpha]$ and (i).

Q.E.D.

Corollary 23. $\mathcal{D}[\alpha]$ is equivalent with each one of the following propositions:

i) There exists a well founded ranked strict ordering relation $<$ with domain $\mathcal{N}_\alpha$, the rank function of which has an image $\geq \Delta$.

ii) There exists a function with domain $\mathcal{N}_\alpha$ which is surjective onto $\Delta$.

$\mathcal{D}[\alpha] \Rightarrow (i)$:

By definition of $\mathcal{D}[\alpha]$, the binary relation $(x = ax) \wedge (y = ay) \wedge (\langle x < y \rangle = \alpha)$ is well founded. By lemma 16(iii), this relation is ranked. We have seen, in the proof of theorem 22, that the image of its rank function is the whole of $\text{On}$.

(i) $\Rightarrow$ (ii) : obvious.

(ii) $\Rightarrow \mathcal{D}[\alpha]$:

Since $\mathcal{D}$ is an ultrafilter, it suffices to show $\neg \mathcal{D}[\neg \alpha]$. But, (ii) and $\mathcal{D}[\neg \alpha]$ contradict lemma 20.

Q.E.D.

Theorem 24.

If $\mathbb{2}$ is non trivial, there exists no set, which is totally ordered by $\epsilon$, the ordinal of which is $\geq \Delta$.

Let $\alpha \in \mathbb{2}$, $\alpha \neq 0, 1$ and $X$ be a set which is totally ordered by $\epsilon$, and equipotent with $\Delta$.

Then, we show that the application $x \mapsto ax$ is an injection from $X$ into $\mathcal{N}_\alpha$:

Indeed, by lemma 16(i), we have $x \epsilon y \rightarrow \langle x < y \rangle = 1$ and, by lemma 5, we have:

$\langle x < y \rangle = 1 \rightarrow \langle ax < ay \rangle = \alpha$. Therefore, if $x, y \in X$ and $x \neq y$, we have, for instance $x \epsilon y$, therefore $\langle ax < ay \rangle = \alpha$ and therefore $ax \neq ay$ since $\alpha \neq 0$.

Thus, there exists a function with domain $\mathcal{N}_\alpha$ which is surjective onto $\Delta$. The same reasoning, applied to $\neg \alpha$ gives the same result for $\neg \alpha$. But this contradicts lemma 20.

Q.E.D.

Remark. Theorem 24 shows that it is impossible to define Von Neumann ordinals in $\mathcal{N}$, with $\epsilon$ instead of $\epsilon$, unless $\mathbb{2}$ is trivial, i.e. the realizability model is, in fact, a forcing model.

The model $M_\mathcal{D}$

For each formula $F[x_1, \ldots, x_n]$ of ZF, we have defined, in the ground model $\mathcal{M}$, an $n$-ary function symbol with values in $\{0, 1\}$, denoted by $\langle F[x_1, \ldots, x_n] \rangle$, by setting, for any individ-
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Proof by recurrence on the length of $F$

If $F$ is atomic, we have $I \models \forall \vec{x} \forall y (\forall y F[\vec{x}, y]) \iff \langle F[\vec{x}, y] \rangle$

and $I \models \forall \vec{x} (\forall \vec{y} \forall y \langle F[\vec{x}, \vec{y}] \rangle)$ since $\langle F[\vec{x}] \rangle$ is identical with $\langle F[\vec{x}] \rangle$.

If $F \equiv F_0 \rightarrow F_1$, the formula $\langle \mathcal{M}_f \models F \rangle \iff \langle \mathcal{D} \langle F \rangle \rangle$ is :

$\langle \mathcal{M}_f \models F \rangle \iff \langle \mathcal{D} \langle F \rangle \rangle$.

Since $\mathcal{D}$ is an ultrafilter, this formula is equivalent with :

$\langle \mathcal{M}_f \models F \rangle \iff \langle \mathcal{D} \langle F \rangle \rangle$, which is a logical consequence of :

$\langle \mathcal{D} \langle F \rangle \rangle$ and $\langle \mathcal{D} \langle F \rangle \rangle$.

Hence the result, by the recurrence hypothesis.

If $F[\vec{x}] \equiv \forall y G[\vec{x}, y]$, let $f_G(\vec{x})$ be the Skolem function of $G$.

Then, we have $\mathcal{M}_f \models \forall y G[\vec{x}, y]$, and therefore :

$I \models \forall \vec{x} \forall y G[\vec{x}, y] \iff \mathcal{M}_f \models G[\vec{x}, f_G(\vec{x})]$.

Therefore, by the recurrence hypothesis, we have :

$I \models \forall \vec{x} \forall y G[\vec{x}, y] \rightarrow \langle G[\vec{x}, f_G(\vec{x})] \rangle$.

Applying Lemma 25(ii), we obtain $I \models \forall \vec{x} \forall y G[\vec{x}, y] \rightarrow \langle \forall \vec{y} \forall y G[\vec{x}, y] \rangle$.

Conversely, by Lemma 25(i), we have $I \models \forall y \langle \forall \vec{y} \forall y G[\vec{x}, y] \rangle$.

Therefore, applying the recurrence hypothesis, we obtain :

$I \models \forall \vec{x} \forall y G[\vec{x}, y] \rightarrow \forall y \mathcal{M}_f \models G[\vec{x}, y]$.

Q.E.D.
Theorem 27. \( \mathcal{M} \) is an elementary extension of the ground model \( \mathcal{M} \).

Let \( F[\bar{a}] \) be a closed formula of ZF, with parameters \( a_1, \ldots, a_n \) in \( \mathcal{M} \).
If \( \mathcal{M} \models F[\bar{a}] \), we have \( \langle F[\bar{a}] \rangle = 1 \) (by definition), and therefore, of course, \( \models \mathcal{D}(F[\bar{a}]) \).
Therefore, by lemma 26, we have \( \models (\mathcal{M} \models F[\bar{a}]). \)
If \( \mathcal{M} \not\models F[\bar{a}] \), then \( \mathcal{M} \models \neg F[\bar{a}] \); therefore, we have \( \models (\mathcal{M} \models \neg F[\bar{a}]) \).
Q.E.D.

Remark. Theorem 27 is, in fact, true for any ultrafilter on \( \mathcal{I} \), with the same proof.

Theorem 28. Let \( \sqsubset \) be a well founded binary relation, defined in the ground model \( \mathcal{M} \). Then the relation \( \mathcal{D}(x \sqsubset y) \) is well founded in the realizability model \( \mathcal{N} \).

Remark. Theorem 28 is an improvement on theorem 8.

Notations. We shall write \( x \sqsubset y \) for \( \langle x \sqsubset y \rangle \in \mathcal{D} \).
Recall that \( x < y \) means \( x \in \text{Cl}(y) \); and that \( x \prec_a y \) means \( \langle x < y \rangle \geq \alpha \), for \( \alpha \in \mathcal{I} \).

We define, in the model \( \mathcal{M} \), a binary relation \( \sqsubset \) on the class \( \{0,1\} \times \mathcal{M} \), by setting, for any \( \alpha, x, a \in \mathcal{M} \):
\[
(a', a) \sqsubset (a, a) \iff (a' < a) \lor (\alpha = a' = 0 \land a' < a) \lor (\alpha = a' = 1 \land a' \not\sqsubset a).
\]
The relation \( \sqsubset \) is the ordered direct sum of the relations \( \sqsubset, < \).
It is easily shown that it is well founded in \( \mathcal{M} \).

The binary function symbol associated with this relation, of domain \( \{0,1\} \times \mathcal{M} \) and values in \( \{0,1\} \), is given by:
\[
\langle (a', a) \sqsubset (a, a) \rangle = (\neg a' \land \alpha) \lor (\neg a' \land \neg a \land (a' < a)) \lor (a' \land a \land (a' \not\sqsubset a)).
\]
This definition gives, in \( \mathcal{N} \), a binary function symbol with arguments in \( \mathcal{I} \times \mathcal{N} \), and values in \( \mathcal{I} \).
By theorem 8, the binary relation \( \langle (a', a) \sqsubset (a, a) \rangle = 1 \) is well founded in \( \mathcal{N} \).

Proof of theorem 28.
Proof by contradiction : we assume that the binary relation \( \sqsubset \) is not well founded.
Thus, there exists \( a_0, A_0 \) such that \( a_0 \in A_0 \) and \( A_0 \) has no minimal \( \epsilon \)-element for \( \sqsubset \).
We define, in \( \mathcal{N} \), the class \( \mathcal{X} \) of ordered pairs \( (a, x) \), such that:
\( (a, x) \in \mathcal{N} \) and \( x \not\prec_a X \) has no minimal \( \epsilon \)-element, neither for \( \sqsubset \) nor for \( \prec_a \).
Therefore, the formula \( \mathcal{X}(a, x) \) is:
\[
a \in \mathcal{I} \land \exists X \forall x \in X \exists \forall u \in X \exists v \in X \exists w \in X \exists u \in \mathcal{I} \land \neg \alpha \land \neg \prec_a \land \neg \sqsubset.
\]
If \( (a, x) \in \mathcal{X} \), then we have \( \mathcal{D}(a) \) : indeed, the set \( X \) is non void and has no minimal \( \epsilon \)-element for \( \prec_a \).
Therefore, we have \( \neg \mathcal{D}(\alpha) \), and thus \( \mathcal{D}(a) \), since \( \mathcal{D} \) is an ultrafilter.
We obtain the desired contradiction by showing that the class \( \mathcal{X} \) is non void and has no minimal element for the binary relation \( \langle (a', x') \sqsubset (a, x) \rangle = 1 \).
The ordered pair \( (1, a_0) \) is in \( \mathcal{X} \) : indeed, we have \( x <_a x \) for every \( x \), and therefore \( a_0 \) has no minimal \( \epsilon \)-element for \( <_a \).
Now let \( (a, a) \) be in \( \mathcal{X} \); we search for \( (a', a') \) in \( \mathcal{X} \) such that \( \langle (a', a') \sqsubset (a, a) \rangle = 1 \).
By hypothesis on \( (a, a) \), there exists \( A \) such that \( a \in A \) and \( A \) has no minimal \( \epsilon \)-element, neither for \( \sqsubset \) nor for \( \prec_a \).
Thus, there exists \( a^0, a^1 \in A \) such that we have \( \mathcal{D}(a^0 \sqsubset a) \) and \( a^1 <_a a \).
Therefore, we have $a$ form a partition of 1 in the Boolean algebra $\mathcal{D}$.

We have $\neg\mathcal{D}(\beta)$; therefore, by definition of $\mathcal{D}$, the relation $<_{\beta}$ is not well founded. Thus, there exists $b, B$ such that $b \in B$ and $B$ has no minimal $\epsilon$-element for $<_{\beta}$. Then, we set:

$$a' = a' a^0 \sqcup (\neg a) a^1 \sqcup b B$$

and $A' = \{x x \sqcup (\neg a) y \sqcup b z ; x, y \in A, z \in B\}$.

Therefore, we have $a' \in A'$, as needed; moreover:

$$\neg a' \wedge a' < a = \neg a,$$ since $\neg a' \geq \neg a$ and $(a' < a) \geq \neg a \wedge (a^1 < a) = \neg a$;

$$a' \wedge a' (a' \sqcup a) = a' \wedge (a^0 \sqcup a) = a'.

By definition of $\langle (a', a') \sqcup (a, a) \rangle$, it follows that $\langle (a', a') \sqcup (a, a) \rangle = \beta \wedge a' = 1$.

It remains to show that $A'$ has no minimal $\epsilon$-element for $\sqsubseteq$ and for $<_{\neg a'}$.

Therefore, let $u \in A'$, thus $u = a' x \sqcup (\neg a) y \sqcup b z$ with $x, y \in A$ and $z \in B$.

By hypothesis on $A, B$, there exists $x', y' \in A, x', y' \sqsubseteq x, y' <_{\neg a} y$ and $z' \in B, z' <_{\beta} z$. Then, if we set $u' = a' x' \sqcup (\neg a) y' \sqcup b z'$, we have $u' \in A'$.

Moreover, we have $\langle u' \sqsubseteq u \rangle \geq a' \wedge (x' \sqsubseteq x)$, and therefore $\mathcal{D}(u' \sqsubseteq u)$, that is $u' \sqsubseteq u$.

Finally, $\langle u' < u \rangle \geq (\neg a \wedge (y' < y)) v (\beta \wedge (z' < z)) = \neg a \wedge \beta = \neg a'$; therefore, we have $u' <_{\neg a'} u$.

Q.E.D.

**Theorem 29.** $\mathcal{M}_\mathcal{D}$ is well founded, and therefore has the same ordinals as $\mathcal{N}_{\epsilon'}$.

We apply theorem 28 to the binary relation $\epsilon$ which is well founded in $\mathcal{M}$. We deduce that the relation $\mathcal{D}(x \in y)$, that is $x \in \mathcal{D} y$, is well founded in $\mathcal{N}$.

Q.E.D.

The relation $\epsilon_\mathcal{D}$ is well founded and extensional, which means that we have, in $\mathcal{N}$:

$$\forall x \forall y (\forall z (z \in \mathcal{D} x \iff z \in \mathcal{D} y) \rightarrow \forall z (x \in \mathcal{D} z \rightarrow y \in \mathcal{D} z))$$

It follows that we can define a collapsing, by means of a function symbol $\Phi$, which is an isomorphism of $(\mathcal{M}_\mathcal{D}, \epsilon_\mathcal{D})$ on a transitive class in the model $\mathcal{N}_\epsilon$ of ZF, which contains the ordinals. This means that we have:

$$\forall x \forall y (\mathcal{D} y \rightarrow \Phi(y) \in \Phi(x)) \wedge \forall x (\forall z \in \Phi(x))(\exists y \in \mathcal{D} x)z = \Phi(y).$$

The definition of $\Phi$ is analogous with that of the rank function already defined for a transitive well founded relation. The details will be given in a later version of this paper.

It follows that:

**Theorem 30.** The realizability model $\mathcal{N}_\epsilon$ contains a transitive class, which contains the ordinals and is an elementary extension of the ground model $\mathcal{M}$.

**Corollary 31.** The class $L^{\mathcal{M}}$ of constructible sets in $\mathcal{M}$ is an elementary submodel of $L^{\mathcal{N}}$.

**References**


Essays on combinatory logic, λ-calculus, and formalism, J.P. Seldin and J.R. Hindley ed.,


