

Coherent Differentiation

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Intro: differentiation and addition

We have learned at school

$$f'(x) = \lim_{\varepsilon \rightarrow 0} \frac{f(x + \varepsilon) - f(x)}{\varepsilon}$$

And later

$f : E \rightarrow F$ (where E and F are, say, Banach spaces)
is differentiable at $x \in E$ if

$$f(x + u) = f(x) + (l \cdot u) + o(\|u\|)$$

where $l : E \rightarrow F$ linear bounded.

And then $l \in \mathcal{L}(E, F)$ is uniquely defined: $l = f'(x)$ is the differential (Jacobian etc) of f at x .

Because $l \cdot u \in o(\|u\|) \Rightarrow l = 0$ when $l \in \mathcal{L}(E, F)$.

Leibniz rule

Take $f : \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$ (sufficiently regular) and define

$$\begin{aligned}g : \mathbb{R} &\rightarrow \mathbb{R} \\x &\mapsto f(x, x)\end{aligned}$$

Then

$$\frac{dg(x)}{dx} = \frac{\partial f(x_1, x)}{\partial x_1}(x) + \frac{\partial f(x, x_2)}{\partial x_2}(x)$$

This generalizes the usual Leibniz rule $(uv)' = u'v + uv'$,
 $(uv)^{(n)} = \sum_{k=0}^n \binom{n}{k} u^{(k)} v^{(n-k)}$ etc.

Differentiation is inherently related to addition.

In the Differential λ -calculus we have a differential application

$$\frac{\Gamma \vdash M : A \Rightarrow B \quad \Gamma \vdash N : A}{\Gamma \vdash DM \cdot N : A \Rightarrow B}$$

and a differential substitution defined by induction on M , such that if $\Gamma, x : A \vdash M : B$ and $\Gamma \vdash N : A$ then

$$\Gamma, x : A \vdash \frac{\partial M}{\partial x} \cdot N : B$$

Differential reduction:

$$D(\lambda x^A M) \cdot N \rightarrow \lambda x^A \left(\frac{\partial M}{\partial x} \cdot N \right)$$

where $\frac{\partial M}{\partial x} \cdot N$ is defined by induction on M .

The most important case in the definition of $\frac{\partial M}{\partial x} \cdot N$ is when $M = (P)Q$:

$$\frac{\partial(P)Q}{\partial x} \cdot N = \left(\frac{\partial P}{\partial x} \cdot N\right)Q + (DP \cdot \left(\frac{\partial Q}{\partial x} \cdot N\right))Q$$

which combines

- the Leibniz Rule because x can occur in P and in Q
- and the Chain Rule because of the application (imagine x occurs only in Q).

Reduction rule:

$$D(\lambda x^A M) \cdot N \rightarrow \lambda x^A \left(\frac{\partial M}{\partial x} \cdot N \right)$$

so to have subject reduction it seems that we need

$$\frac{\Gamma \vdash M_0 : A \quad \Gamma \vdash M_1 : A}{\Gamma \vdash M_0 + M_1 : A}$$

allowing to add **any two terms of the same type**.

Consequence: non-determinism

If we have for instance a type o of booleans with

$$\overline{\Gamma \vdash \mathbf{t} : o} \quad \overline{\Gamma \vdash \mathbf{f} : o}$$

then we *must* accept $\mathbf{t} + \mathbf{f}$ as a valid term, with

$$\Gamma \vdash \mathbf{t} + \mathbf{f} : o$$

meaning that the language is essentially non-deterministic.

In the semantics

So far, the categorical models \mathbf{C} of the differential λ -calculus were (left-)additive categories.

Given $f, g \in \mathbf{C}(A, B)$, there is a morphism $f + g \in \mathbf{C}(A, B)$.

$\rightsquigarrow \mathbf{C}$ is enriched over commutative monoids.

Coherent Differentiation

This is not a fatality!

Fact

*Of course addition is required, but there is a (categorical, and then syntactical) way of controlling it, **without giving up determinism**.*

The possibility of such a theory appears in ...

... probabilistic coherence spaces (PCS)

A PCS is a pair $X = (|X|, PX)$ where $|X|$ is a set and $PX \subseteq (\mathbb{R}_{\geq 0})^{|X|}$ satisfying some closure properties.

- PX is convex,
- downwards closed,
- closed under lubs of monotonic ω -chains
- + a technical condition to avoid ∞ coeffs.

They are a model of (probabilistic) λ -calculi, LL etc, but **not of their differential extensions** by lack of additivity.

Derivatives in PCSs

In the associated category \mathbf{Pcoh}_1 , 1 is an object such that $|1| = \{*\}$, $P1 = [0, 1]$ and a morphism $f \in \mathbf{Pcoh}_1(1, 1)$ is a power series defining a function $[0, 1] \rightarrow [0, 1]$:

$$f(x) = \sum_{n=0}^{\infty} a_n x^n \quad \text{with} \quad \forall n \ a_n \in \mathbb{R}_{\geq 0} \quad \text{and} \quad \sum_{n=0}^{\infty} a_n \leq 1$$

so that $f'(x) = \sum_{n=0}^{\infty} (n+1)a_{n+1}x^n$ has no reason to be a function $[0, 1] \rightarrow [0, 1]$.

Example

$f(x) = 1 - \sqrt{1-x}$, then $f'(x) = 1/(2\sqrt{1-x})$ is unbounded on $[x, 1)$.

However

Fact

If $x, u \in [0, 1]$ and $x + u \in [0, 1]$ then we have

$$f(x) + f'(x)u \leq f(x + u) \in [0, 1]$$

because this sum is the beginning of the Taylor expansion which holds in this model:

$$f(x + u) = \sum_{n=0}^{\infty} \frac{1}{n!} f^{(n)}(x) u^n$$

and all coefficients are ≥ 0 . For any $f \in \mathbf{Pcoh}_!(X, Y)$.

In a PCS, some sums are allowed...

Convex combinations: if $x, y \in PX$ then $\frac{1}{3}x + \frac{2}{3}y \in PX$.

Some **non convex** sums are also possible, for instance in the cartesian product $X \& Y$, we have

$$(x, 0) + (0, y) = (x, y) \in P(X \& Y) = PX \times PY$$

if $x \in PX$ and $y \in PY$.

Other **non convex** allowed sums come from differentiation:

$$f(x) + f'(x) \cdot u \in PY$$

if $x, u \in PX$ are such that $x + y \in PX$ and $f : PX \rightarrow PY$ is an “analytic function”, that is a morphism $X \rightarrow Y$ in the Kleisli category **Pcoh**₁.

... and some sums are forbidden!

For instance

$$P(1 \oplus 1) = \{(x_0, x_1) \in \mathbb{R}_{\geq 0} \mid x_0 + x_1 \leq 1\}$$

in this object of booleans,

$$\mathbf{t} = (1, 0), \mathbf{f} = (0, 1) \in P(1 \oplus 1) \quad \text{and} \quad \mathbf{t} + \mathbf{f} \notin P(1 \oplus 1).$$

or simply $1 \in [0, 1]$ and $1 + 1 \notin [0, 1]$.

Fundamental observation

There is a functor $\mathbf{S} : \mathbf{Pcoh} \rightarrow \mathbf{Pcoh}$ which maps an object X to an object $\mathbf{S}X$ such that

$$P(\mathbf{S}X) = \{(x, u) \in PX^2 \mid x + u \in PX\}.$$

For instance $P(\mathbf{S}1) = \{(x, u) \in [0, 1] \mid x + u \leq 1\}$.

We base our axiomatization on the existence of such a functor.

Summable categories

Definition (pre-summable category)

A *pre-summable* category is a tuple

$$(\mathcal{L}, \mathbf{S}, \pi_0, \pi_1, \sigma)$$

Where

- \mathcal{L} is a category enriched over pointed sets (and the distinguished morphism is always denoted 0);
- $\mathbf{S} : \mathcal{L} \rightarrow \mathcal{L}$ is a functor which preserves the enrichment ($\mathbf{S}0 = 0$);
- and $\pi_0, \pi_1, \sigma : \mathbf{S}X \rightarrow X$ are natural transformations such that π_0 and π_1 are jointly monic.

If $f_0, f_1 \in \mathcal{L}(X, \mathbf{S}Y)$ satisfy $\pi_j f_0 = \pi_j f_1$ for $j = 0, 1$ then $f_0 = f_1$.

Intuition

- $\mathbf{S}X$ is the objects of pairs $(x_0, x_1) \in X \times X$ such that $x_0 + x_1$ is well defined and $\in X$;
- $\pi_j : \mathbf{S}X \rightarrow X$ are the projections, $\pi_j(x_0, x_1) = x_j$;
- and $\sigma : \mathbf{S}X \rightarrow X$ maps (x_0, x_1) to $x_0 + x_1$.

Some terminology

We assume to have such a structure $(\mathcal{L}, \mathbf{S}, \pi_0, \pi_1, \sigma)$

Definition (summability, witness, sum)

We say that $f_0, f_1 \in \mathcal{L}(X, Y)$ are **summable** if there is $h \in \mathcal{L}(X, \mathbf{S}Y)$ such that $\pi_j h = f_j$ for $j = 0, 1$.

Fact: when such an h exists it is unique (π_0, π_1 are jointly monic), we set $\langle f_0, f_1 \rangle_{\mathbf{S}} = h$, it is the **witness of summability** of f_0 and f_1 .

And then we set $f_0 + f_1 = \sigma \langle f_0, f_1 \rangle_{\mathbf{S}}$, the **sum** of f_0 and f_1 .

Some simple observations

- π_0, π_1 are summable with $\langle \pi_0, \pi_1 \rangle_{\mathbf{s}} = \text{Id}_{\mathbf{s}X}$ and $\pi_0 + \pi_1 = \sigma$.
- If $f_0, f_1 \in \mathcal{L}(X, Y)$ are summable and $l \in \mathcal{L}(U, X)$ and $r \in \mathcal{L}(Y, V)$ then $r f_0 l, r f_1 l$ are summable with

$$\begin{aligned}\langle r f_0 l, r f_1 l \rangle_{\mathbf{s}} &= \mathbf{S}r \langle f_0, f_1 \rangle_{\mathbf{s}} l \\ r f_0 l + r f_1 l &= r (f_0 + f_1) l\end{aligned}$$

by naturality of π_0, π_1 and σ .

Remark (main tool)

Use the fact that π_0, π_1 are jointly monic.

We introduce a few axioms to make this “partial addition” behave as expected.

Commutativity

Axiom (Commutativity)

π_1, π_0 are summable and $\pi_1 + \pi_0 = \sigma$.

Fact (consequences)

$\langle \pi_1, \pi_0 \rangle_{\mathbf{S}} \in \mathcal{L}(\mathbf{S}X, \mathbf{S}X)$ is an involution.

If $f_0, f_1 \in \mathcal{L}(X, Y)$ are summable then f_1, f_0 are summable with $f_1 + f_0 = f_0 + f_1$.

Intuitively: $\langle \pi_1, \pi_0 \rangle_{\mathbf{S}}(x_0, x_1) = (x_1, x_0)$.

Neutrality

Axiom (Neutrality)

For any $f \in \mathcal{L}(X, Y)$, f and 0 are summable and $f + 0 = f$.

In particular we have two injections

$$\iota_0 = \langle \text{Id}_X, 0 \rangle_{\mathbf{S}}, \iota_1 = \langle 0, \text{Id}_X \rangle_{\mathbf{S}} \in \mathcal{L}(X, \mathbf{S}X)$$

Intuitively $\iota_0(x) = (x, 0)$ and $\iota_1(x) = (0, x)$.

Witness

Associativity is more tricky. We split the condition in two pieces.

Axiom (Witness)

Let $f_{ij} \in \mathcal{L}(X, Y)$ for $i, j \in \{0, 1\}$ be 4 morphisms such that

- f_{j0}, f_{j1} are summable for $j = 0, 1$
- and $f_{00} + f_{01}, f_{10} + f_{11}$ are summable

then $\langle f_{00}, f_{01} \rangle_{\mathbf{S}}, \langle f_{10}, f_{11} \rangle_{\mathbf{S}}$ are summable.

So there is a witness for this summability:

$$\langle \langle f_{00}, f_{01} \rangle_{\mathbf{S}}, \langle f_{10}, f_{11} \rangle_{\mathbf{S}} \rangle_{\mathbf{S}} \in \mathcal{L}(X, \mathbf{S}^2 Y).$$

The canonical flip

Fact

There is exactly one morphism $c \in \mathcal{L}(\mathbf{S}^2X, \mathbf{S}^2X)$ such that

$$\forall i, j \in \{0, 1\} \quad \pi_i \pi_j c = \pi_j \pi_i .$$

$$c = \langle \langle \pi_0 \pi_0, \pi_0 \pi_1 \rangle \mathbf{s}, \langle \pi_1 \pi_0, \pi_0 \pi_0 \rangle \mathbf{s} \rangle \mathbf{s}$$

exists by the previous axioms.

Fact

$$c^2 = \text{Id}_{\mathbf{S}^2X} .$$

Intuitively $c((x_{00}, x_{01}), (x_{10}, x_{11})) = ((x_{00}, x_{10}), (x_{01}, x_{11}))$.

Associativity

Axiom (Associativity)

The following diagram commutes

$$\begin{array}{ccc} \mathbf{S}^2 X & \xrightarrow{c} & \mathbf{S}^2 X \\ & \searrow \sigma_{\mathbf{S}X} & \swarrow \mathbf{S}\sigma_X \\ & \mathbf{S}X & \end{array}$$

Remark (Intuition)

The sum of witnesses is performed componentwise:

$$\langle x_{00}, x_{01} \rangle_{\mathbf{S}} + \langle x_{10}, x_{11} \rangle_{\mathbf{S}} = \langle x_{00} + x_{10}, x_{01} + x_{11} \rangle_{\mathbf{S}}$$

Fact (consequence)

If $f_{ij} \in \mathcal{L}(X, Y)$ for $i, j \in \{0, 1\}$ are such that

- f_{j0}, f_{j1} are summable for $j = 0, 1$
- and $f_{00} + f_{01}, f_{10} + f_{11}$ are summable

then

- f_{0j}, f_{1j} are summable for $j = 0, 1$
- and $f_{00} + f_{10}, f_{01} + f_{11}$ are summable

and $(f_{00} + f_{01}) + (f_{10} + f_{11}) = (f_{00} + f_{10}) + (f_{01} + f_{11})$.

Associativity follows taking $f_{10} = 0$.

Partially additive category

The category \mathcal{L} becomes a partially additive category in the sense of partial monoids.

Remark

Partially additive categories do not suffice for our goal: the functor **S** will be crucial for differentiation!

\mathbf{S} is a monad

We have already $\zeta = \iota_0 = \langle \text{Id}_X, 0 \rangle_{\mathbf{S}} \in \mathcal{L}(X, \mathbf{S}X)$.

Using the axioms we also have

$$\theta = \langle \pi_0 \pi_0, \pi_1 \pi_0 + \pi_0 \pi_1 \rangle_{\mathbf{S}} \in \mathcal{L}(\mathbf{S}^2X, \mathbf{S}X).$$

Fact

$(\mathbf{S}, \zeta, \theta)$ is a monad on \mathcal{L} .

Intuitively

$$\begin{aligned} \theta_X : \mathbf{S}^2X &\rightarrow \mathbf{S}X \\ ((x_{00}, x_{01}), (x_{10}, x_{11})) &\mapsto (x_{00}, x_{10} + x_{01}) \end{aligned}$$

Notice that **we forget x_{11}** .

When \mathcal{L} is an SMC

Distributivity

In all the situations we have in mind, \mathcal{L} is a symmetric monoidal category with tensor product \otimes and tensor unit 1.

In that case one expects \otimes to distribute over $+$, when defined. This requires an additional

Axiom (Distributivity)

$$0 \otimes g = 0$$

and

if f_0, f_1 are summable then

- $f_0 \otimes g, f_1 \otimes g$ are summable
- and $f_0 \otimes g + f_1 \otimes g = (f_0 + f_1) \otimes g$.

Strength

In particular $\pi_0 \otimes \text{Id}_Y, \pi_1 \otimes \text{Id}_Y \in \mathcal{L}(\mathbf{S}X \otimes Y, X \otimes Y)$ are summable so we have strengths

$$\varphi_{X,Y}^0 = \langle \pi_0 \otimes \text{Id}_Y, \pi_1 \otimes \text{Id}_Y \rangle_{\mathbf{S}} \in \mathcal{L}(\mathbf{S}X \otimes Y, \mathbf{S}(X \otimes Y))$$

$$\varphi_{X,Y}^1 = \langle \text{Id}_X \otimes \pi_0, \text{Id}_X \otimes \pi_1 \rangle_{\mathbf{S}} \in \mathcal{L}(X \otimes \mathbf{S}Y, \mathbf{S}(X \otimes Y))$$

turn \mathbf{S} into a commutative monad.

Intuitively

$$\varphi_{X,Y}^0((x_0, x_1) \otimes y) = (x_0 \otimes y, x_1 \otimes y)$$

$$\varphi_{X,Y}^1(x \otimes (y_0, y_1)) = (x \otimes y_0, x \otimes y_1)$$

Commutativity of the monad

We have actually something stronger:

$$\begin{array}{ccccc} \mathbf{S}X \otimes \mathbf{S}Y & \xrightarrow{\varphi_{X,\mathbf{S}Y}^0} & \mathbf{S}(X \otimes \mathbf{S}Y) & \xrightarrow{\mathbf{S}\varphi_{X,Y}^1} & \mathbf{S}^2(X \otimes Y) \\ \varphi_{\mathbf{S}X,Y}^1 \downarrow & & & & \downarrow c_{X \otimes Y} \\ \mathbf{S}(\mathbf{S}X \otimes Y) & \xrightarrow{\mathbf{S}\varphi_{X,Y}^0} & & & \mathbf{S}^2(X \otimes Y) \end{array}$$

Intuitively

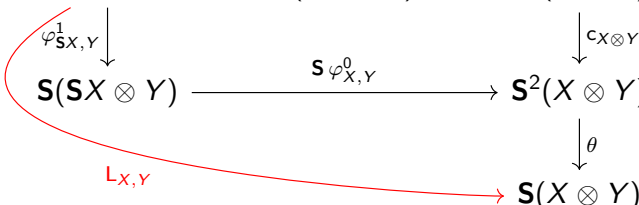
$$(x_0, x_1) \otimes (y_0, y_1) \mapsto ((x_0 \otimes y_0, x_0 \otimes y_1), (x_1 \otimes y_0, x_1 \otimes y_1))$$

$$(x_0, x_1) \otimes (y_0, y_1) \mapsto ((x_0 \otimes y_0, x_1 \otimes y_0), (x_0 \otimes y_1, x_1 \otimes y_1))$$

Induced symmetric monoidal structure

We then have

$$\begin{array}{ccccc}
 \mathbf{S}X \otimes \mathbf{S}Y & \xrightarrow{\varphi_{X,\mathbf{S}Y}^0} & \mathbf{S}(X \otimes \mathbf{S}Y) & \xrightarrow{\mathbf{S}\varphi_{X,Y}^1} & \mathbf{S}^2(X \otimes Y) \\
 \downarrow \varphi_{\mathbf{S}X,Y}^1 & & & & \downarrow c_{X \otimes Y} \\
 \mathbf{S}(\mathbf{S}X \otimes Y) & \xrightarrow{\mathbf{S}\varphi_{X,Y}^0} & & \longrightarrow & \mathbf{S}^2(X \otimes Y) \\
 & & & & \downarrow \theta \\
 & & & & \mathbf{S}(X \otimes Y)
 \end{array}$$

$L_{X,Y}$


Intuitively

$$L_{X,Y} : ((x_0, x_1) \otimes (y_0, y_1)) \mapsto (x_0 \otimes y_0, x_0 \otimes y_1 + x_1 \otimes y_0)$$

Differential structure

As in differential LL, we consider differentiation as a structure of the exponential.

So we assume moreover that

- \mathcal{L} is cartesian (\top : terminal object, $X_0 \& X_1$: product, $\text{pr}_i \in \mathcal{L}(X_0 \& X_1, X_i)$, $\langle f_0, f_1 \rangle \in \mathcal{L}(Y, X_0 \& X_1)$ if $f_i \in \mathcal{L}(Y, X_i)$).
- \mathcal{L} is equipped with a resource modality ($!-, \text{der}, \text{dig}, m^0, m^2$)

$$\text{der}_X \in \mathcal{L}(!X, X) \quad \text{dig}_X \in \mathcal{L}(!X, !!X) \quad \text{comonad structure}$$
$$m^0 \in \mathcal{L}(1, !\top) \quad m^2_{X,Y} \in \mathcal{L}(!X \otimes !Y, !(X \& Y))$$

Seelye isos, strong sym. monoidality

Preservation of products

We need a further property about **S**.

Axiom (Product)

The functor **S** preserves cartesian products, more precisely:

$$\langle \mathbf{S}pr_0, \mathbf{S}pr_1 \rangle \in \mathcal{L}(\mathbf{S}(X_0 \& X_1), \mathbf{S}X_0 \& \mathbf{S}X_1)$$

is an iso.

This holds in all the LL-based examples we have in mind, because in these examples **S** is a right adjoint.

The differentiation operator

In this setting (resource category with a summability structure), a *differential structure* is a natural transformation

$$\partial_X \in \mathcal{L}(!\mathbf{S}X, \mathbf{S}!X)$$

satisfying some properties.

Remark (main idea)

Given $f \in \mathcal{L}_!(X, Y)$, this will allow to define

$$\mathbf{D}f = (\mathbf{S}f) \partial_X \in \mathcal{L}_!(\mathbf{S}X, \mathbf{S}Y)$$

which will (intuitively) be the map $(x, u) \mapsto (f(x), f'(x) \cdot u)$.

We list the conditions to be satisfied by ∂_X

Second derivative: intuition

Let $f \in \mathcal{L}_1(X, Y)$, we have $\mathbf{D}f \in \mathcal{L}_1(\mathbf{D}X, \mathbf{D}Y)$

$$\mathbf{D}f(x, u) = \left(f(x), \frac{df(x)}{dx} \cdot u \right)$$

We can apply \mathbf{D} to $\mathbf{D}f$, we get

$$\mathbf{D}^2 f((x, u), (y, v)) = \left(\mathbf{D}f(x, u), \frac{d\mathbf{D}f(x, u)}{d(x, u)} \cdot (y, v) \right)$$

Remember $\mathbf{D}f(x, u) = (f(x), f'(x) \cdot u)$.

By standard rules of calculus:

$$\frac{d\mathbf{D}f(x, u)}{d(x, u)} \cdot (y, v) = \frac{\partial \mathbf{D}f(x, u)}{\partial x} \cdot y + \frac{\partial \mathbf{D}f(x, u)}{\partial u} \cdot v$$

$$\begin{aligned} \frac{\partial \mathbf{D}f(x, u)}{\partial x} \cdot y &= \frac{\partial}{\partial x} (f(x), f'(x) \cdot u) \cdot y \\ &= (f'(x) \cdot y, f''(x) \cdot (u, y)) \end{aligned}$$

$$\begin{aligned} \frac{\partial \mathbf{D}f(x, u)}{\partial u} \cdot v &= \frac{\partial}{\partial u} (f(x), f'(x) \cdot u) \cdot v \\ &= (0, f'(x) \cdot v) \end{aligned}$$

Finally we have, intuitively

$$\mathbf{D}^2 f((x, u), (y, v)) = ((f(x), f'(x) \cdot u), \\ (f'(x) \cdot y, f''(x) \cdot (u, y) + f'(x) \cdot v))$$

Notice that in the first 3 components, we have only 1st order derivatives.

Distributive law

Axiom (Chain Rule + Linearity)

∂ is a distributive law between the monad \mathbf{S} and the comonad $!_-$ in the following sense.

$$\begin{array}{ccc}
 !\mathbf{S}X & \xrightarrow{\partial_X} & \mathbf{S}!X \\
 \searrow \text{ders}_X & & \downarrow \mathbf{S} \text{der}_X \\
 & & \mathbf{S}X
 \end{array}
 \qquad
 \begin{array}{ccc}
 !\mathbf{S}X & \xrightarrow{\partial_X} & \mathbf{S}!X \\
 \text{digs}_X \downarrow & & \downarrow \mathbf{S} \text{digs}_X \\
 !!\mathbf{S}X & \xrightarrow{!\partial_X} & !\mathbf{S}!X \xrightarrow{\partial_{!X}} & \mathbf{S}!!X
 \end{array}$$

$$\begin{array}{ccc}
 !\mathbf{S}X & \xrightarrow{\partial_X} & \mathbf{S}!X \\
 \swarrow !\zeta_X & & \uparrow \zeta_{!X} \\
 & & !X
 \end{array}
 \qquad
 \begin{array}{ccc}
 !\mathbf{S}X & \xrightarrow{\partial_X} & \mathbf{S}!X \\
 !\theta_X \uparrow & & \uparrow \theta_{!X} \\
 !\mathbf{S}^2X & \xrightarrow{\partial_{\mathbf{S}X}} & \mathbf{S}!\mathbf{S}X \xrightarrow{\mathbf{S}\partial_X} & \mathbf{S}^2!X
 \end{array}$$

See John Power and Hiroshi Watanabe, *Combining a monad and a comonad*, TCS 2002 for this kind of dist. law.

Intuition for the dist. law

The **first two diagrams** allow to define a functor

$$\begin{aligned} \mathbf{D} : \mathcal{L}_! &\rightarrow \mathcal{L}_! \\ X &\mapsto \mathbf{S}X \\ (f : !X \rightarrow Y) &\mapsto ((\mathbf{S}f) \partial_X : !\mathbf{S}X \rightarrow \mathbf{S}Y) \end{aligned}$$

Intuitively, and in probabilistic coherence spaces for instance:

- $f \in \mathcal{L}_!(X, Y)$ means that f is an analytic function $X \rightarrow Y$
- $\mathbf{D}f \in \mathbf{Pcoh}_!(X, Y)$ is the $(x, u) \mapsto (f(x), f'(x) \cdot u)$

so this functoriality means that the chain rule holds.

And that the differential of a linear morphism is the morphism itself: $\mathbf{D}(f \text{ der}_X) = (\mathbf{S}f) \text{ der}_{\mathbf{S}X}$ for $f \in \mathcal{L}(X, Y)$.

The two next diagrams allow to lift the monad $(\mathbf{S}, \zeta, \theta)$ to $\mathcal{L}_!$.

For $\theta_X \in \mathcal{L}_!(\mathbf{S}^2X, \mathbf{S}Y) = \mathcal{L}(!\mathbf{S}^2X, \mathbf{S}Y)$: we take $\theta_X \text{ der}_X$.

These diagrams allow to prove that θ is a natural transformation on $\mathcal{L}_!$. If $f \in \mathcal{L}_!(X, Y)$:

$$\begin{array}{ccc} \mathbf{D}^2X & \xrightarrow{\theta_X} & \mathbf{D}X \\ \mathbf{D}^2f \downarrow & & \downarrow \mathbf{D}f \\ \mathbf{D}^2Y & \xrightarrow{\theta_Y} & \mathbf{D}Y \end{array}$$

And similarly ζ is natural in $\mathcal{L}_!$.

Intuition: linearity of the differential

Remember:

$$\theta_X((x_0, u_0), (x_1, u_1)) = (x_0, u_0 + x_1)$$

$$\mathbf{D}^2 f((x_0, u_0), (x_1, u_1)) = ((f(x_0), f'(x_0) \cdot u_0), \\ (f'(x_0) \cdot x_1, f''(x_0) \cdot (u_0, x_1) + f'(x_0) \cdot u_1))$$

The commutation means:

$$\mathbf{D}f(x_0, u_0 + x_1) = (f(x_0), f'(x_0) \cdot u_0 + f'(x_0) \cdot x_1)$$

that is $f'(x_0) \cdot (u_0 + x_1) = f'(x_0) \cdot u_0 + f'(x_0) \cdot x_1$.

Naturality of ζ in \mathcal{L}_1 : $f'(x) \cdot 0 = 0$.

Locality

To represent one of the differential situation we are interested in, this distributive law has to satisfy additional axioms: *Locality*, *Leibniz* and *Schwarz*.

Axiom (Locality)

$$\begin{array}{ccc} !S X & \xrightarrow{\partial_X} & S !X \\ \searrow \! \pi_0 & & \swarrow \pi_0 \\ & !X & \end{array}$$

Only for π_0 , not for π_1 !

Intuition

Again we use π_0 for $\pi_0 \operatorname{der}_X \in \mathcal{L}_!(\mathbf{D}X, X)$.

The diagram means that π_0 is natural in $\mathcal{L}_!$. If $f \in \mathcal{L}_!(X, Y)$:

$$\begin{array}{ccc} \mathbf{D}X & \xrightarrow{\pi_0} & X \\ \mathbf{D}f \downarrow & & \downarrow f \\ \mathbf{D}Y & \xrightarrow{\pi_0} & Y \end{array}$$

This corresponds to the intuition that

$$\mathbf{D}f(x, u) = (f(x), f'(x) \cdot u)$$

Remark

$\pi_1 \in \mathcal{L}_!(\mathbf{D}X, X)$ also exists but is fundamentally **not natural** in $\mathcal{L}_!$ (of course π_1 is natural in \mathcal{L}).

Leibniz

Is expressed as a “monoidality” condition (to simplify we assume $\mathbf{S}(X \& Y) = \mathbf{S}X \& \mathbf{S}Y$)

Axiom (Leibniz)

$$\begin{array}{ccc} \mathbf{!S}X \otimes \mathbf{!S}Y & \xrightarrow{\partial_X \otimes \partial_Y} & \mathbf{S!}X \otimes \mathbf{S!}Y \xrightarrow{L_{!X, !Y}} \mathbf{S}(!X \otimes !Y) \\ m_{\mathbf{S}X, \mathbf{S}Y}^2 \downarrow & & \downarrow \mathbf{S}m_{X, Y}^2 \\ \mathbf{!S}(X \& Y) & \xrightarrow{\partial_{X \& Y}} & \mathbf{S}!(X \& Y) \end{array}$$

+ a “0-ary version”.

Intuition

Given $f \in \mathcal{L}_!(X \& Y, Z)$, this commutation gives us an expression for $\mathbf{D}f$ in terms of the two differentials ∂_X and ∂_Y .

Given $((x, y), (u, v)) \in \mathbf{S}(X \& Y)$,
that is, $(x, u) \in \mathbf{S}X$ and $(y, v) \in \mathbf{S}Y$,

$$\frac{df(x, y)}{d(x, y)} \cdot (u, v) = \frac{\partial f(x, y)}{\partial x} \cdot u + \frac{\partial f(x, y)}{\partial y} \cdot v$$

In the diagram, $+$ is implemented by $L_{!X, !Y}$.

Axiom (Schwarz)

$$\begin{array}{ccccc}
 \mathbf{!S^2X} & \xrightarrow{\partial_{\mathbf{s}X}} & \mathbf{S!SX} & \xrightarrow{\mathbf{S}\partial_X} & \mathbf{S^2!X} \\
 \downarrow \mathbf{!c_X} & & & & \downarrow \mathbf{c_{!X}} \\
 \mathbf{!S^2X} & \xrightarrow{\partial_{\mathbf{s}X}} & \mathbf{S!SX} & \xrightarrow{\mathbf{S}\partial_X} & \mathbf{S^2!X}
 \end{array}$$

If $f \in \mathcal{L}_1(X, Y)$ then

$$\mathbf{D}^2 f = (\mathbf{S}^2 f) (\mathbf{S} \partial_X) \partial_{\mathbf{S}X}$$

so this diagram means that c is natural in \mathcal{L}_1 :

$$\begin{array}{ccc} \mathbf{D}^2 X & \xrightarrow{\mathbf{D}^2 f} & \mathbf{D}^2 Y \\ c_X \downarrow & & \downarrow c_Y \\ \mathbf{D}^2 X & \xrightarrow{\mathbf{D}^2 f} & \mathbf{D}^2 Y \end{array}$$

where we use also c_X for $c_X \text{ der}_{\mathbf{S}^2 X}$.

Intuition

Remember that

$$c_X((x_0, u_0), (x_1, u_1)) = ((x_0, x_1), (u_0, u_1))$$

$$\mathbf{D}^2 f((x_0, u_0), (x_1, u_1)) = ((f(x_0), f'(x_0) \cdot u_0), \\ (f'(x_0) \cdot x_1, f''(x_0) \cdot (u_0, x_1) + f'(x_0) \cdot u_1))$$

so this naturality means that

$$\mathbf{D}^2 f((x_0, x_1), (u_0, u_1)) = ((f(x_0), f'(x_0) \cdot x_1), \\ (f'(x_0) \cdot u_0, f''(x_0) \cdot (u_0, x_1) + f'(x_0) \cdot u_1))$$

$$((f(x_0), f'(x_0) \cdot x_1), (f'(x_0) \cdot u_0, f''(x_0) \cdot (x_1, u_0) + f'(x_0) \cdot u_1)) \\ = ((f(x_0), f'(x_0) \cdot x_1), (f'(x_0) \cdot u_0, f''(x_0) \cdot (u_0, x_1) + f'(x_0) \cdot u_1))$$

So taking $u_1 = 0$ we get

$$f''(x_0) \cdot (x_1, u_0) = f''(x_0) \cdot (u_0, x_1)$$

which is the crucial property that the second derivative is a **symmetric** bilinear function, often called *Schwarz lemma*.

Coherent Differentiation (II)

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Short recap

Summability structure

- \mathcal{L} is a category with 0-morphisms
- $\mathbf{S} : \mathcal{L} \rightarrow \mathcal{L}$ is a 0-preserving functor
- $\pi_0, \pi_1, \sigma : \mathbf{S}X \rightarrow X$ are natural transformations
- π_0, π_1 are jointly monic
- $f_0, f_1 \in \mathcal{L}(X, Y)$ are summable if there is $\langle f_0, f_1 \rangle_{\mathbf{S}} \in \mathcal{L}(X, \mathbf{S}Y)$ with $\pi_i \langle f_0, f_1 \rangle_{\mathbf{S}} = f_i$, and then $f_0 + f_1 = \sigma \langle f_0, f_1 \rangle_{\mathbf{S}}$.

+ axioms to turn $\mathcal{L}(X, Y)$ into a partial commutative monoid.

In particular $c \in \mathcal{L}(\mathbf{S}^2X, \mathbf{S}^2X)$ the standard flip with

$$\pi_i \pi_j c = \pi_j \pi_i.$$

\mathbf{S} inherits a monad structure $(\mathbf{S}, \zeta, \theta)$.

Differentiation

\mathcal{L} is assumed to be a resource category (cartesian SMC with a resource comonad aka. exponential, with Seely strong monoidality).

The differential structure is a natural transformation

$\partial_X \in \mathcal{L}(!\mathbf{S}X, \mathbf{S}!X)$ which satisfies some further commutations:

- it is a distributive law between the monad \mathbf{S} and the comonad $!_:$ Chain Rule and Linearity (of the derivative)
- Locality
- Leibniz
- Schwarz.

Then one defines the **Differentiation Functor** $\mathbf{D} : \mathcal{L}_! \rightarrow \mathcal{L}_!$ by

$\mathbf{D}X = \mathbf{S}X$ and if $f \in \mathcal{L}_!(X, Y) = \mathcal{L}_!(!X, Y)$ then

$\mathbf{D}f = (\mathbf{S}f) \partial_X \in \mathcal{L}_!(\mathbf{D}X, \mathbf{D}Y)$.

Canonical structure

A special, very common, case

We assume that \mathcal{L} is monoidal **closed** (convenient though not strictly necessary) so that the functor

$$- \otimes I$$

where $I = 1 \ \& \ 1$, has a right adjoint S_I .

$$S_I X = (I \multimap X)$$

$$S_I f = (I \multimap f) \in \mathcal{L}(I \multimap X, I \multimap Y)$$

for $f \in \mathcal{L}(X, Y)$.

Remark

We still assume that \mathcal{L} has zero-morphisms.

Two natural questions

Remark

In (probabilistic) coherence spaces, \mathbf{S} is defined exactly in that way.

- When does this definition give rise to a summability structure?
- What does the differential structure boil down to in this setting?

We have three morphisms

$$\bar{\pi}_0 = \langle \text{Id}, 0 \rangle \in \mathcal{L}(1, I)$$

$$\bar{\pi}_1 = \langle 0, \text{Id} \rangle \in \mathcal{L}(1, I)$$

$$\Delta = \langle \text{Id}, \text{Id} \rangle \in \mathcal{L}(I, I)$$

which induce natural transformations $\pi_0, \pi_1, \sigma \in \mathcal{L}(\mathbf{S}_I X, X)$ by “precomposition”.

For instance π_0 is

$$(I \multimap X) \xrightarrow{\sim} (I \multimap X) \otimes 1 \xrightarrow{\text{Id} \otimes \bar{\pi}_0} (I \multimap X) \otimes I \xrightarrow{\text{ev}} X$$

Summability as a property

Definition

\mathcal{L} is **canonically summable** if $(\mathbf{S}_I, \pi_0, \pi_1, \sigma)$ defined in that way are a summability structure.

Remark (a property of \mathcal{L} , not a structure)

This is a **property** of \mathcal{L} , not an additional structure on \mathcal{L} .

In particular we need $\bar{\pi}_0, \bar{\pi}_1$ to be jointly epic.

What do summability and sums become?

Remember that $\mathcal{L}(X, Y) \simeq \mathcal{L}(1, X \multimap Y)$.

Fact

$x_0, x_1 \in \mathcal{L}(1, X)$ are summable if there is $h \in \mathcal{L}(1, X)$ such that

$$x_i = h \bar{\pi}_i$$

and then $x_0 + x_1 = h \Delta \in \mathcal{L}(1, X)$.

Canonical Witness Axiom

If $f_0, f_1 \in \mathcal{L}(I, X)$ are such that $f_0 \Delta, f_1 \Delta \in \mathcal{L}(I, X)$ are summable, then so are f_0, f_1 . That is, up to $\mathcal{L}(I, X) \simeq \mathcal{L}(1, I \multimap X)$:

if $f_0, f_1, f \in \mathcal{L}(I, X)$ are such that

$$f_i \Delta = f \bar{\pi}_i \text{ for } i = 0, 1$$

then there is $h \in \mathcal{L}(I \otimes I, X)$ such that

$$f_i \lambda = h(\bar{\pi}_i \otimes I) \in \mathcal{L}(1 \otimes I, X)$$

where λ is the can. isom. $1 \otimes I \rightarrow I$.

Remark

Then $f \rho = h(I \otimes \Delta)$. Because $\bar{\pi}_0, \bar{\pi}_1$ are jointly epic.

Theorem

If $\bar{\pi}_0, \bar{\pi}_1$ are jointly epic, then $(\mathbf{S}_I, \pi_0, \pi_1, \sigma)$ (as defined above) is a summability structure on \mathcal{L} iff the Canonical Witness Axiom holds.

I is a commutative comonoid

Thanks to the axioms we can define

$$\tilde{L} \in \mathcal{L}(I, I \otimes I)$$

uniquely characterized by

$$\tilde{L}\pi_0 = \pi_0 \otimes \pi_0 \text{ and } \tilde{L}\pi_1 = \pi_0 \otimes \pi_1 + \pi_1 \otimes \pi_0$$

Fact

$(I, \text{pr}_0 \in \mathcal{L}(I, 1), \tilde{L})$ is a commutative comonoid in \mathcal{L} .

$\text{pr}_0 \in \mathcal{L}(I = (1 \& 1), 1)$ is the first projection.

The commutative monad structure of \mathbf{S}_I

We have seen that \mathbf{S}_I has a structure of commutative monad.

Fact

The monad $(\mathbf{S}_I, \zeta, \theta)$ is induced by the commutative comonoid structure (pr_0, \tilde{L}) of I .

For instance $\theta = \text{cur } f : (I \multimap (I \multimap X)) \rightarrow (I \multimap X)$ where f is

$$\begin{array}{ccc} (I \multimap (I \multimap X)) \otimes I & \xrightarrow{\text{Id} \otimes \tilde{L}} & (I \multimap (I \multimap X)) \otimes I \otimes I \\ & & \downarrow \text{ev} \otimes \text{Id} \\ X & \xleftarrow{\text{ev}} & (I \multimap X) \otimes I \end{array}$$

Differentiation as a $!$ -coalgebra (canonical case)

!_ and its coalgebras

We assume that \mathcal{L} is a cartesian resource category (cartesian product $\&$, exponential comonad $!_$, Seelye isos etc).

A **!-coalgebra** structure on $X \in \mathcal{L}$ is a $d \in \mathcal{L}(X, !X)$ such that

$$\begin{array}{ccc} X & \xrightarrow{d} & !X \\ & \searrow \text{Id} & \downarrow \text{der}_X \\ & & X \end{array} \qquad \begin{array}{ccc} X & \xrightarrow{d} & !X \\ d \downarrow & & \downarrow \text{dig}_X \\ !X & \xrightarrow{!d} & !!X \end{array}$$

These colgebras form the **Eilenberg-Moore** category $\mathcal{L}^!$ where

$$f \in \mathcal{L}^!((X, d), (Y, e)) \text{ if } \begin{array}{ccc} X & \xrightarrow{d} & !X \\ f \downarrow & & \downarrow !f \\ Y & \xrightarrow{e} & !Y \end{array} \text{ in } \mathcal{L}.$$

$\mathcal{L}^!$ is cartesian

Due to the fact that \mathcal{L} is a resource category (\otimes , $\&$, Seely isos):

Fact

$\mathcal{L}^!$ is cartesian, with terminal object

$$(1, \mu^0 : 1 \rightarrow !1)$$

and the cartesian product of $(P_0, d_0), (P_1, d_1) \in \mathcal{L}^!$ is $(P_0 \otimes P_1, \mu^2(d_0 \otimes d_1))$

$$P_0 \otimes P_1 \xrightarrow{d_0 \otimes d_1} !P_0 \otimes !P_1 \xrightarrow{\mu^2} !(P_0 \otimes P_1)$$

Projection pr_0^\otimes (and similarly for pr_1^\otimes):

$$P_0 \otimes P_1 \xrightarrow{d_0 \otimes d_1} !P_0 \otimes !P_1 \xrightarrow{\text{der}_{P_0} \otimes !0} P_0 \otimes !\top \simeq P_0$$

Uses the lax symmetric monoidality structure (μ^0, μ^2) of $! \dots$

Chain Rule and coalgebra

Fact

There is a bijective correspondence between

- the $!$ -coalgebra structures on I
- and the distributive laws between \mathbf{S}_I and $!_-$ in the sense of the *Chain Rule*:

$$\begin{array}{ccc} !SX & \xrightarrow{\partial_X} & S!X \\ & \searrow \text{ders}_X & \downarrow \mathbf{S} \text{ der}_X \\ & & SX \end{array}$$

$$\begin{array}{ccccc} !SX & \xrightarrow{\partial_X} & S!X & & \\ \text{digs}_X \downarrow & & \downarrow \mathbf{S} \text{ dig}_X & & \\ !!SX & \xrightarrow{! \partial_X} & !S!X & \xrightarrow{\partial_{!X}} & S!!X \end{array}$$

Coalgebra \rightsquigarrow Chain Rule

Suppose we are given $\delta \in \mathcal{L}(I, !I)$, then for any object X we can define $\partial_X = \text{cur } f \in \mathcal{L}(!\mathbf{S}_I X = !(I \multimap X), \mathbf{S}_I !X = (I \multimap !X))$ where f is

$$!(I \multimap X) \otimes I \xrightarrow{\text{Id} \otimes \delta} !(I \multimap X) \otimes !I \xrightarrow{\mu^2} !((I \multimap X) \otimes I) \xrightarrow{!ev} !X$$

$\mu_{X,Y}^2 \in \mathcal{L}(!X \otimes !Y, !(X \otimes Y))$ is the lax-monoidality structure of $!_-$ wrt. \otimes .

Chain Rule \rightsquigarrow Coalgebra

Conversely assume we are given $\partial_X \in \mathcal{L}(!\mathbf{S}_1 X, \mathbf{S}_1 !X)$ for each $X \in \mathcal{L}$, we have in particular, taking $X = I$:

$$\begin{array}{ccccc}
 I & \xrightarrow{\lambda_I^{-1}} & 1 \otimes I & \xrightarrow{\mu^0 \otimes \text{Id}} & !1 \otimes I & \xrightarrow{!(\text{cur } \lambda_I) \otimes \text{Id}} & !(I \multimap I) \otimes I \\
 & & & & & \searrow \partial_I \otimes \text{Id} & \\
 & & & & & & (I \multimap !I) \otimes I & \xrightarrow{\text{ev}} & !I
 \end{array}$$

where $\mu^0 \in \mathcal{L}(1, !1)$ is the “unit” of the lax-monoidality and $\lambda_I \in \mathcal{L}(1 \otimes I, I)$ (the canonical iso).

A natural question

So assume we are given a coalgebra structure $\delta \in \mathcal{L}(I, !I)$.

What conditions must satisfy δ for ensuring that the corresponding distributive law $(\partial_X)_{X \in \mathcal{L}}$ satisfies the additional conditions

- **Linearity** (second part of the dist. law)
- Local
- Leibniz
- Schwarz?

The answer is surprisingly simple.

Linearity and Leibniz

Linearity **and** Leibniz boil down to

$$\begin{array}{ccc}
 | & \xrightarrow{\delta} & !! \\
 \text{pr}_0 \downarrow & & \downarrow !\text{pr}_0 \\
 1 & \xrightarrow{\mu^0} & !1
 \end{array}
 \qquad
 \begin{array}{ccc}
 | & \xrightarrow{\delta} & !! \\
 \tilde{\text{L}} \downarrow & & \downarrow !\tilde{\text{L}} \\
 | \otimes | & \xrightarrow{\delta \otimes \delta} & !! \otimes !! \xrightarrow{\mu^2} & !(| \otimes |)
 \end{array}$$

that is

$$\begin{array}{c}
 \text{term. obj.} \\
 \text{pr}_0 \in \mathcal{L}^!((|, \delta), \overbrace{(1, \mu^0)}) \\
 \tilde{\text{L}} \in \mathcal{L}^!((|, \delta), \underbrace{(|, \delta) \otimes (|, \delta)}_{\text{cart. prod.}})
 \end{array}$$

comonoid from the coalgebra

This means that we have

$$\begin{array}{ccc}
 I & \xrightarrow{\delta} & !! \\
 \text{pr}_0 \downarrow & & \downarrow !0 \\
 1 & \xleftarrow{(m^0)^{-1}} & !\top
 \end{array}
 \qquad
 \begin{array}{ccc}
 I & \xrightarrow{\delta} & !! \\
 \tilde{L} \downarrow & & \downarrow \Delta \\
 I \otimes I & \xleftarrow{\text{der}_1 \otimes \text{der}_1} & !! \otimes !! \xleftarrow{(m^2)^{-1}} & !(I \& I)
 \end{array}$$

because $\mathcal{L}^!$ is cartesian.

Remark

As a consequence, a canonically summable resource category where $!_-$ is the free exponential (roughly speaking, a Lafont category which is canonically summable) has exactly one differential structure (in our sense).

Related to a result of Blute, Cockett, Lemay and Seely (in additive resource categories).

Locality corresponds to

$$\begin{array}{ccc} \mathbb{1} & \xrightarrow{\delta} & \mathbb{!}\mathbb{1} \\ \iota_0 \uparrow & & \uparrow \mathbb{!}\iota_0 \\ \mathbb{1} & \xrightarrow{\mu^0} & \mathbb{!}\mathbb{1} \end{array}$$

that is $\iota_0 \in \mathcal{L}^!((\mathbb{1}, \mu^0), (\mathbb{1}, \delta))$.

And Schwarz straightforwardly holds.

Remark: the Kleisli category of \mathbf{S}_1

It turns out to be exactly the same thing as the category $\mathcal{L}[(I, \delta)]$ of **free comodules** of the coalgebra (I, δ) .

Theorem (Girard)

If \mathcal{L} is a model of LL then $\mathcal{L}[(I, \delta)]$ is a model of LL. Very likely conjecture: it is also a summable differential model of LL.

The objects of $\mathcal{L}[(I, \delta)]$ are those of \mathcal{L} .

$f \in \mathcal{L}[(I, \delta)](X, Y)$ if $f = (f_0, f_1) \in \mathcal{L}(X, Y)$ is a summable pair of morphisms. Composition:

$$(g_0, g_1) (f_0, f_1) = (g_0 f_0, g_0 f_1 + g_1 f_0).$$

Intuition: “ $g_1 f_1 = 0$ ”, $\mathcal{L}[(I, \delta)]$ is a kind of infinitesimal extension of \mathcal{L} .

To summarize

In the canonical case, for a closed resource category \mathcal{L} :

- 1 summability boils down to the Canonical Witness Axiom about $I = 1 \ \& \ 1$ (+ the fact that $\bar{\pi}_0, \bar{\pi}_1$ are jointly epic);
- 2 and the differential structure boils down to a coalgebra structure on I

such that the morphisms $\text{pr}_0 \in \mathcal{L}(I, 1)$, $\iota_0 \in \mathcal{L}(1, I)$ and $\tilde{L} \in \mathcal{L}(I, I \otimes I)$ are coalgebra morphisms.

Remember that these 3 morphisms arise from the summability assumptions.

Concrete instance I: Coherence Spaces

A coherence space is

$$E = (|E|, \circlearrowright_E)$$

where $|E|$ is a set and \circlearrowright_E is a binary symmetric and reflexive relation on $|E|$.

The domain of cliques:

$$\text{Cl}(E) = \{x \subseteq |E| \mid \forall a, a' \in x \ a \circlearrowright_E a'\}$$

ordered by \subseteq .

Morphisms

- $|E \multimap F| = |E| \times |F|$
- $(a, b) \circ_{E \multimap F} (a', b')$ if

$$a \circ_E a' \Rightarrow (b \circ_F b' \text{ and } b = b' \Rightarrow a = a')$$

And then

$$\mathbf{Coh}(E, F) = \text{Cl}(E \multimap F)$$

Some notations for **Coh**

- Identity: $\text{Id}_E = \{(a, a) \mid a \in |E|\}$
- Composition: if $s \in \mathbf{Coh}(E, F)$ and $t \in \mathbf{Coh}(F, G)$ then

$$ts = \{(a, c) \in |E| \times |G| \mid \exists b \in |F| (a, b) \in s \text{ and } (b, c) \in t\} \\ \in \mathbf{Coh}(E, G)$$

- Application to a clique: if $s \in \mathbf{Coh}(E, F)$ and $x \in \text{Cl}(E)$ then $s \cdot x = \{b \in |F| \mid a \in x \text{ and } (a, b) \in s\} \in \text{Cl}(F)$.

Coh is cartesian

- Terminal object $\top = (\emptyset, \emptyset)$.
- Cartesian product $|E_0 \& E_1| = \{0\} \times |E_0| \cup \{1\} \times |E_1|$
 $(i, a) \circ_{E_0 \& E_1} (j, b)$ if $i = j \Rightarrow a \circ_{E_i} b$.
- The projections are

$$\text{pr}_i = \{((i, a), a) \mid i \in \{0, 1\} \text{ and } a \in |E_i|\} \in \mathbf{Coh}(E_0 \& E_1, E_i).$$

If $t_i \in \mathbf{Coh}(F, E_i)$ then

$$\langle t_0, t_1 \rangle = \{(b, (i, a)) \mid i \in \{0, 1\} \text{ and } (b, a) \in t_i\} \\ \in \mathbf{Coh}(F, E_0 \& E_1).$$

Remark

$\text{Cl}(\top) = \{\emptyset\}$ and $\text{Cl}(E_0) \times \text{Cl}(E_1) \simeq \text{Cl}(E_0 \& E_1)$ by

$$(x_0, x_1) \mapsto \{0\} \times x_0 \cup \{1\} \times x_1.$$

Coh is monoidal closed

- Unit $1 = (\{*\}, =)$.
- Tensor product $|E_0 \otimes E_1| = |E_0| \times |E_1|$ and $(a_0, a_1) \supset_{E_0 \otimes E_1} (a'_0, a'_1)$ if $a_i \supset_{E_i} a'_i$ for $i = 0, 1$.
- If $t_i \in \mathbf{Coh}(E_i, F_i)$ for $i = 0, 1$ then

$$t_0 \otimes t_1 = \{((a_0, a_1), (b_0, b_1)) \mid (a_i, b_i) \in t_i \text{ for } i = 0, 1\} \\ \in \mathbf{Coh}(E_0 \otimes E_1, F_0 \otimes F_1).$$

Monoidal closedness:

$$\mathbf{Coh}(G \otimes E, F) \simeq \mathbf{Coh}(G, E \multimap F).$$

Coh as a resource category

- $!E$ = the set of all finite multisets $m = [a_1, \dots, a_n]$ with $a_i \in E$ and $\forall i, j \ a_i \supset_E a_j$. It is a **uniform** exponential.
- $m \supset_{!E} m'$ if $\forall a \in m, a' \in m' \ m \supset_E m'$.
- And if $t \in \mathbf{Coh}(E, F)$ then

$$\begin{aligned} !t = \{ & ([a_1, \dots, a_n], [b_1, \dots, b_n]) \mid \\ & n \in \mathbb{N}, [a_1, \dots, a_n] \in !E \\ & \text{and } (a_i, b_i) \in t \text{ for } i = 1, \dots, n\} \\ & \in \mathbf{Coh}(!E, !F). \end{aligned}$$

Remark

This is the free exponential. There is another one where $!E$ is made of **sets** instead of multisets; it is not compatible with the differential structure.

Coh is canonically summable

- **Coh** has 0-morphisms: $0 = \emptyset \in \mathbf{Coh}(E, F)$.
- $I = 1$ & 1 so that $|I| = \{0, 1\}$ and $0 \subset_I 1$.
- The injections $\bar{\pi}_i = \{(*, i)\} \in \mathbf{Coh}(1, I)$ are jointly epic.

Remark

$s \in \mathbf{Coh}(I, E)$ is fully determined by the pair

$$s \cdot \{0\}, s \cdot \{1\} \in \text{Cl}(E)$$

such that

$$s \cdot \{0\} \cap s \cdot \{1\} = \emptyset.$$

The Can. Witness Axiom holds in **Coh**

Let $t_0, t_1, t \in \mathbf{Coh}(I, E)$ such that

$$t_i \Delta = t \bar{\pi}_i \text{ for } i = 0, 1.$$

This means $t_i \cdot \{0, 1\} = t \cdot \{i\}$ for $i = 0, 1$. That is:

$$t_0 \cdot \{0, 1\} \cup t_1 \cdot \{0, 1\} \in \text{Cl}(E) \text{ and } t_0 \cdot \{0, 1\} \cap t_1 \cdot \{0, 1\} = \emptyset.$$

Then let $s = \{((i, j), a) \mid (i, a) \in t_j\} \subseteq |I \otimes I \multimap E|$, we have

$$s \in \mathbf{Coh}(I \otimes I, E).$$

The functor $\mathbf{S}_I : \mathbf{Coh} \rightarrow \mathbf{Coh}$ is given by

$$\mathbf{S}_I E = (I \multimap E)$$

so that $|\mathbf{S}_I E| = \{0, 1\} \times |E|$ with

$$(i, a) \circ_{\mathbf{S}_I E} (i', a') \text{ if } a \circ_E a' \text{ and } i \neq i' \Rightarrow a \neq a'.$$

Hence

$$\text{Cl}(\mathbf{S}_I E) \simeq \{(x_0, x_1) \in \text{Cl}(E)^2 \mid x_0 \cup x_1 \in \text{Cl}(E) \text{ and } x_0 \cap x_1 = \emptyset\}.$$

Remark

Two cliques x_0, x_1 of E are summable if $x_0 \cup x_1 \in \text{Cl}(E)$ and $x_0 \cap x_1 = \emptyset$. In that case we use $x_0 + x_1$ for $x_0 \cup x_1$.

The commutative comonoid structure of I is given by

$$\text{pr}_0 = \{(0, *)\} \in \mathbf{Coh}(I, 1)$$

$$\tilde{I} = \{(0, (0, 0)), (1, (1, 0)), (1, (0, 1))\} \in \mathbf{Coh}(I, I \otimes I).$$

Remember it induces the monad structure $\zeta_E \in \mathbf{Coh}(E, \mathbf{S}_I E)$ and $\theta_E \in \mathbf{Coh}(\mathbf{S}_I^2 E, \mathbf{S}_I E)$.

As expected

$$\begin{aligned} \theta_E : \mathbf{S}_I^2 E &\rightarrow \mathbf{S}_I E \\ ((x_{00}, x_{01}), (x_{10}, x_{11})) &\mapsto (x_{00}, x_{10} + x_{01}) \end{aligned}$$

up to

$$\text{Cl}(\mathbf{S}_I^2 E) \simeq \{((x_{00}, x_{01}), (x_{10}, x_{11})) \mid x_{00} + x_{01} + x_{10} + x_{11} \in \text{Cl}(E)\}.$$

The differential structure of **Coh**

We define $\delta \subseteq |I \multimap !!|$:

$$\delta = \{(0, n[0]) \mid n \in \mathbb{N}\} \cup \{(1, n[0] + [1]) \mid n \in \mathbb{N}\}$$

where $n[a] = [\overbrace{a, \dots, a}^{n \times}]$. It is easy to check that $\delta \in \mathbf{Coh}(I, !!)$.

δ is a coalgebra

The main thing to check is

$$\begin{array}{ccc} | & \xrightarrow{\delta} & !! \\ \delta \downarrow & & \downarrow !\delta \\ !! & \xrightarrow{\text{dig}_E} & !!! \end{array}$$

that is, given $i \in \{0, 1\}$ and $M \in \mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(\{0, 1\}))$,

$$(i, M) \in !\delta \delta \Leftrightarrow (i, M) \in \text{dig}_! \delta$$

where

$$\text{dig}_E = \{(m, [m_1, \dots, m_k]) \in |E| \times !!E \mid m = m_1 + \dots + m_k\}.$$

main case

The main case is when $i = 1$.

$(1, M) \in !\delta$ means $\exists k \in \mathbb{N}$ such that

$$(k[0] + [1], M) \in !\delta$$

that is:

$$M = [m_1, \dots, m_{k+1}] \text{ with } (0, m_i) \in \delta \text{ for } i = 1, \dots, k \\ \text{and } (1, m_{k+1}) \in \delta$$

that is: $\exists k \in \mathbb{N} \exists n_1, \dots, n_{k+1} \in \mathbb{N}$

$$M = [n_1[0], \dots, n_k[0], n_{k+1}[0] + [1]]$$

And $(1, M) \in \text{dig}_l \delta$ means $\exists k \in \mathbb{N}$ such that

$$(k[0] + [1], M) \in \text{dig}_l$$

that is:

$$M = [m_1, \dots, m_l] \text{ with } m_1 + \dots + m_l = k[0] + [1]$$

that is: $\exists l \in \mathbb{N}^+ \exists n_1, \dots, n_l \in \mathbb{N}$

$$M = [n_1[0], \dots, n_{l-1}[0], n_l[0] + [1]]$$

The diagram commutes!

The differential distributive law

Remember that δ induces a distributive law

$\partial_E = \text{cur } u \in \mathbf{Coh}(!\mathbf{S}_1 E, \mathbf{S}_1 !E)$ where

$$u : !(I \multimap E) \otimes I \rightarrow !E$$

is

$$!(I \multimap E) \otimes I \xrightarrow{\text{Id} \otimes \delta} !(I \multimap E) \otimes !I \xrightarrow{\mu^2} !((I \multimap E) \otimes I) \xrightarrow{!ev} !E$$

Notice that $\mu_{E,F}^2 \in \mathbf{Coh}(!E \otimes !F, !(E \otimes F))$ is

$$\mu_{E,F}^2 = \{(((a_1, \dots, a_n), [b_1, \dots, b_n]), [(a_1, b_1), \dots, (a_n, b_n)]) \mid n \in \mathbb{N}, [a_1, \dots, a_n] \in !E \text{ and } [b_1, \dots, b_n] \in !F\}$$

$$!(I \multimap E) \otimes !I \xrightarrow{\mu^2} !((I \multimap E) \otimes I) \xrightarrow{!ev} !E$$

is

$$\{(((i_1, a_1), \dots, (i_k, a_k)), [i_1, \dots, i_k]), [a_1, \dots, a_k]) \mid k \in \mathbb{N}, i_1, \dots, i_k \in \{0, 1\} \text{ and } [(i_1, a_1), \dots, (i_k, a_k)] \in !(I \multimap E)\}$$

and $[(i_1, a_1), \dots, (i_k, a_k)] \in !(I \multimap E)$ means that

$$\forall j, j' \ a_j \supset_E a_{j'} \text{ and } j \neq j' \Rightarrow a_j \neq a_{j'}.$$

$$!(I \multimap E) \otimes I \xrightarrow{\text{Id} \otimes \delta} !(I \multimap E) \otimes !I$$

is

$$\begin{aligned} & \{((p, 0), (p, n[0])) \mid n \in \mathbb{N}, p \in \text{Cl}(I \multimap E)\} \\ & \cup \{((p, 1), (p, n[0] + [1])) \mid n \in \mathbb{N}, p \in \text{Cl}(I \multimap E)\} \end{aligned}$$

so $u =$

$$!(I \multimap E) \otimes I \xrightarrow{\text{Id} \otimes \delta} !(I \multimap E) \otimes !I \xrightarrow{\mu^2} !((I \multimap E) \otimes I) \xrightarrow{!ev} !E$$

is

$$u = \{(((0, a_1), \dots, (0, a_k)), 0), [a_1, \dots, a_k]) \mid \\ k \in \mathbb{N} \text{ and } [a_1, \dots, a_k] \in !E\} \\ \cup \{(((0, a_1), \dots, (0, a_k), (1, a_{k+1})), 1), [a_1, \dots, a_{k+1}]) \mid \\ k \in \mathbb{N} \text{ and } [a_1, \dots, a_{k+1}] \in !E \text{ and } a_{k+1} \notin \{a_1, \dots, a_k\}\}$$

Expression of ∂_E

$$\begin{aligned} \partial_E = & \{ (((0, a_1), \dots, (0, a_k)), (0, [a_1, \dots, a_k]) \mid \\ & k \in \mathbb{N} \text{ and } [a_1, \dots, a_k] \in |!E| \} \\ \cup & \{ (((0, a_1), \dots, (0, a_k), (1, a_{k+1})), (1, [a_1, \dots, a_{k+1}]) \mid \\ & k \in \mathbb{N} \text{ and } [a_1, \dots, a_{k+1}] \in |!E| \text{ and } a_{k+1} \notin \{a_1, \dots, a_k\} \} \\ \in & \mathbf{Coh}(!(| \multimap E), | \multimap !E). \end{aligned}$$

The Kleisli category $\mathbf{Coh}_!$

Object: those of \mathbf{Coh} and $\mathbf{Coh}_!(E, F) = \mathbf{Coh}(!E, F)$.

$s \in \mathbf{Coh}_!(E, F)$ induces a **stable function**

$$\hat{s} : \text{Cl}(E) \rightarrow \text{Cl}(F)$$

$$x \mapsto \{b \in |F| \mid \exists m \in \mathcal{M}_{\text{fin}}(x) (m, b) \in s\}$$

Remark

Different s 's can induce the same stable function: \hat{s} forgets about the multiplicities in multisets.

If $s_1 = \{([a], b)\}$ and $s_2 = \{([a, a], b)\}$ then $\hat{s}_1 = \hat{s}_2$.

Differentiation on $\mathbf{Coh}_!$

Given $t \in \mathbf{Coh}_!(E, F) = \mathbf{Coh}(!E, F)$, remember that

$$\mathbf{D}t = (\mathbf{S}t) \partial_E \in \mathbf{Coh}(!\mathbf{S}_!E = (I \multimap E), \mathbf{S}_!E = (I \multimap !F)).$$

Notice that

$$\mathbf{S}t = \{((i, m), (i, b)) \mid i \in \{0, 1\} \text{ and } (m, b) \in t\}.$$

So

$$\begin{aligned} \mathbf{D}t = & \{([(0, a_1), \dots, (0, a_k)], (0, b)) \mid ([a_1, \dots, a_k], b) \in t\} \\ & \cup \{([(0, a_1), \dots, (0, a_k), (1, a_{k+1})], (1, b)) \mid \\ & ([a_1, \dots, a_k, a_{k+1}], b) \in t \text{ and } a_{k+1} \notin \{a_1, \dots, a_k\}\} \end{aligned}$$

The stable derivative

Remember that

$$\text{Cl}(\mathbf{S}_1 E) \simeq \{(x, u) \mid x \cup u \in \text{Cl}(E) \text{ and } x \cap u = \emptyset\}.$$

In that way we get the stable function

$$\begin{aligned} \widehat{\mathbf{D}}t : \text{Cl}(\mathbf{S}_1 E) &\rightarrow \text{Cl}(\mathbf{S}_1 F) \\ (x, u) &\mapsto (\widehat{t}(x), t'(x) \cdot u) \end{aligned}$$

where

$$t'(x) \cdot u = \{b \in |F| \mid \exists m \in \mathcal{M}_{\text{fin}}(x), a \in u (m + [a], b) \in t\}$$

Remark

If $t_i = \{(i[a], a)\}$ for $i = 1, 2$ we get

$$t'_1(\emptyset) \cdot \{a\} = \{a\}$$

$$t'_2(\emptyset) \cdot \{a\} = \emptyset$$

whereas $\widehat{t}_1 = \widehat{t}_2$. The derivative is not associated with the stable function itself.

In some sense this derivative "does not see multiplicities". This can be remedied using non-uniform coherence spaces.

Coherent Differentiation (III)

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The joint “epicness” axiom is necessary

Remember that we required \mathcal{L} to satisfy the following.

The morphisms $\bar{\pi}_0, \bar{\pi}_1 : 1 \rightarrow I = 1 \& 1$ are jointly epic, that is: if $f_0, f_1 : I \rightarrow X$ satisfy

$$f_0 \bar{\pi}_i = f_1 \bar{\pi}_i \quad \text{for } i = 0, 1$$

then $f_0 = f_1$.

This is not always true.

The category of pointed sets

Pointed set: a set X together with a distinguished $0_X \in X$.

Morphisms: functions $f : X \rightarrow Y$ s.t. $f(0_X) = 0_Y$.

This category is cartesian:

$X \& Y = X \times Y$ with $0_{X\&Y} = (0_X, 0_Y)$.

And monoidal closed

$$X \otimes Y = \{(x, y) \in X \times Y \mid x = 0_X \Leftrightarrow y = 0_Y\}$$

and $0_{X \otimes Y} = (0_X, 0_Y)$ (smash product). The \otimes -unit is $1 = \{*, 0_1\}$.

Remark

It is even a resource category: set $!X = \{(0, 0)\} \cup \{1\} \times X$, and $0_{!X} = (0, 0)$.

Then the injections $\bar{\pi}_0, \bar{\pi}_1 : 1 \rightarrow I = 1 \& 1$ are given by

$$\bar{\pi}_0(0_1) = (0_1, 0_1) \quad \bar{\pi}_0(*) = (*, 0_1)$$

$$\bar{\pi}_1(0_1) = (0_1, 0_1) \quad \bar{\pi}_1(*) = (0_1, *)$$

If $f, g : 1 \& 1 \rightarrow Y$ satisfy $f \bar{\pi}_i = g \bar{\pi}_i$ for $i = 0, 1$ we can still have $f(*, *) \neq g(*, *)$.

The CWA is necessary

Canonical Witness Axiom

Reminder

Joint epicness of $\bar{\pi}_0, \bar{\pi}_1$ and CWA are the only conditions we need in the canonical case to get a summability structure.

If $f_0, f_1 \in \mathcal{L}(I, X)$ are such that $f_0 \Delta, f_1 \Delta \in \mathcal{L}(I, X)$ are summable, then so are f_0, f_1 . That is, up to $\mathcal{L}(I, X) \simeq \mathcal{L}(1, I \multimap X)$:

if $f_0, f_1, f \in \mathcal{L}(I, X)$ are such that

$$f_i \Delta = f \bar{\pi}_i : 1 \rightarrow X \text{ for } i = 0, 1$$

then there is $h \in \mathcal{L}(I \otimes I, X)$ such that

$$f_i \lambda = h(\bar{\pi}_i \otimes I) \in \mathcal{L}(1 \otimes I, X)$$

where λ is the can. isom. $1 \otimes I \rightarrow I$.

Normed vector spaces

The CWA do not always hold.

Let \mathcal{N} be the category

- whose objects are the finite-dimensional \mathbb{R} -vector spaces V equipped with a norm $\|\cdot\|_V$
- and a morphism $f : V \rightarrow W$ is a linear map such that $\forall v \in V \ \|f(v)\|_W \leq \|v\|_V$, that is $\|f\| \leq 1$,

where

$$\|f\| = \sup_{\|v\|_V \leq 1} \|f(v)\|_W.$$

\mathcal{N} is cartesian with $\|(u, v)\|_{V \& W} = \max(\|u\|, \|v\|)$ for $(u, v) \in V \& W = V \times W$.

\mathcal{N} is an SMCC with $\|v \otimes w\|_{V \otimes W} = \|v\| \|w\|$ for $v \in V$ and $w \in W$.

The unit of \otimes is $1 = \mathbb{R}$ with $\|r\| = |r|$.

$V \multimap W$ is the space of all linear maps $V \rightarrow W$ with the norm $\|f\|_{V \multimap W} = \|f\|$ already defined.

Joint epicness axiom **holds** in \mathcal{N} .

Then the functor $\mathbf{S}_1 V$ (induced by I) is given by

$$\mathbf{S}_1 V = (I \multimap V) = V \times V \text{ and}$$

$$\|(u, v)\|_{\mathbf{S}_1 V} = \sup_{a, b \in [-1, 1]} \|au + bv\|_V$$

So $u, v \in V$ are summable if $\forall a, b \in [-1, 1] \quad \|au + bv\|_V \leq 1$.

In \mathbb{R} :

- $-1/2$ and $1/2$ are summable since $\|(-1/2, 1/2)\|_{\mathbf{S}_1 \mathbb{R}} = 1$
- $-1/2 + 1/2 = 0$ and 1 are summable
- but $1/2$ and 1 are not summable.

\rightsquigarrow the CWA **does not hold** in \mathcal{N} .

Remark

CWA expresses not only associativity of (partial) $+$ but also some form of **positivity** of the elements of $\mathcal{L}(X, Y)$.

Recap of the differential structure

Assume that \mathcal{L} is a canonically summable resource category, that is:

- $\bar{\pi}_0, \bar{\pi}_1 \in \mathcal{L}(1, I = 1 \ \& \ 1)$ are jointly epic
- and the CWA holds.

Remember that I has a commutative comonoid structure given by

$$\text{pr}_0 : I \rightarrow 1 \qquad \tilde{L} : I \rightarrow I \otimes I$$

with

$$\tilde{L} \bar{\pi}_0 = \bar{\pi}_0 \otimes \bar{\pi}_0 \qquad \tilde{L} \bar{\pi}_1 = \bar{\pi}_0 \otimes \bar{\pi}_1 + \bar{\pi}_1 \otimes \bar{\pi}_0$$

Remember that $+$ is just a notation for a composition with $\Delta = \langle \text{Id}_1, \text{Id}_1 \rangle : 1 \rightarrow I$.

Differential structure

In this setting, a differential structure is a $!$ -coalgebra structure $\delta \in \mathcal{L}(I, !I)$ such that

- $\text{pr}_0 \in \mathcal{L}^!(I, \delta), (1, \mu^0)$
- $\tilde{L} \in \mathcal{L}^!(I, \delta), (I, \delta) \otimes (I, \delta) = (I \otimes I, \mu^2(\delta \otimes \delta))$
- $\bar{\pi}_0 \in \mathcal{L}^!(1, \mu^0), (I, \delta)$.

Theorem

If \mathcal{L} is a Lafont resource category which is canonically summable, then \mathcal{L} has exactly one differential structure.

Idea: δ is uniquely determined by (pr_0, \tilde{L}) .

Lafont resource category: for each $X \in \mathcal{L}$, $!X$ is the free commutative comonoid “cogenerated” by X .

Coherence spaces

A coherence space is

$$E = (|E|, \circlearrowright_E)$$

where $|E|$ is a set and \circlearrowright_E is a binary symmetric and reflexive relation on $|E|$.

The domain of cliques:

$$\text{Cl}(E) = \{x \subseteq |E| \mid \forall a, a' \in |E| \ a \circlearrowright_E a'\}$$

ordered by \subseteq , it is a cpo.

Morphisms

- $|E \multimap F| = |E| \times |F|$
- $(a, b) \circ_{E \multimap F} (a', b')$ if

$$a \circ_E a' \Rightarrow (b \circ_F b' \text{ and } b = b' \Rightarrow a = a')$$

And then

$$\mathbf{Coh}(E, F) = \text{Cl}(E \multimap F)$$

Some notations for **Coh**

- Identity: $\text{Id}_E = \{(a, a) \mid a \in |E|\}$
- Composition: if $s \in \mathbf{Coh}(E, F)$ and $t \in \mathbf{Coh}(F, G)$ then

$$ts = \{(a, c) \in |E| \times |G| \mid \exists b \in |F| (a, b) \in s \text{ and } (b, c) \in t\} \\ \in \mathbf{Coh}(E, G)$$

- Application to a clique: if $s \in \mathbf{Coh}(E, F)$ and $x \in \text{Cl}(E)$ then $s \cdot x = \{b \in |F| \mid a \in x \text{ and } (a, b) \in s\} \in \text{Cl}(F)$.

Coh is cartesian

- Terminal object $\top = (\emptyset, \emptyset)$.
- Cartesian product $|E_0 \& E_1| = \{0\} \times |E_0| \cup \{1\} \times |E_1|$
 $(i, a) \circ_{E_0 \& E_1} (j, b)$ if $i = j \Rightarrow a \circ_{E_i} b$.
- The projections are

$$\text{pr}_i = \{((i, a), a) \mid i \in \{0, 1\} \text{ and } a \in |E_i|\} \in \mathbf{Coh}(E_0 \& E_1, E_i).$$

If $t_i \in \mathbf{Coh}(F, E_i)$ then

$$\langle t_0, t_1 \rangle = \{(b, (i, a)) \mid i \in \{0, 1\} \text{ and } (b, a) \in t_i\} \\ \in \mathbf{Coh}(F, E_0 \& E_1).$$

Remark

$\text{Cl}(\top) = \{\emptyset\}$ and $\text{Cl}(E_0) \times \text{Cl}(E_1) \simeq \text{Cl}(E_0 \& E_1)$ by

$$(x_0, x_1) \mapsto \{0\} \times x_0 \cup \{1\} \times x_1.$$

Coh is monoidal closed

- Unit $1 = (\{*\}, =)$.
- Tensor product $|E_0 \otimes E_1| = |E_0| \times |E_1|$ and $(a_0, a_1) \supset_{E_0 \otimes E_1} (a'_0, a'_1)$ if $a_i \supset_{E_i} a'_i$ for $i = 0, 1$.
- If $t_i \in \mathbf{Coh}(E_i, F_i)$ for $i = 0, 1$ then

$$t_0 \otimes t_1 = \{((a_0, a_1), (b_0, b_1)) \mid (a_i, b_i) \in t_i \text{ for } i = 0, 1\} \\ \in \mathbf{Coh}(E_0 \otimes E_1, F_0 \otimes F_1).$$

Monoidal closedness:

$$\mathbf{Coh}(G \otimes E, F) \simeq \mathbf{Coh}(G, E \multimap F).$$

Coh as a resource category

- $!E$ = the set of all finite multisets $m = [a_1, \dots, a_n]$ with $a_i \in E$ and $\forall i, j \ a_i \supset_E a_j$. It is a **uniform** exponential.
- $m \supset_{!E} m'$ if $\forall a \in m, a' \in m' \ m \supset_E m'$.
- And if $t \in \mathbf{Coh}(E, F)$ then

$$\begin{aligned} !t = \{ & ([a_1, \dots, a_n], [b_1, \dots, b_n]) \mid \\ & n \in \mathbb{N}, [a_1, \dots, a_n] \in !E \\ & \text{and } (a_i, b_i) \in t \text{ for } i = 1, \dots, n\} \\ & \in \mathbf{Coh}(!E, !F). \end{aligned}$$

Remark

This is the free exponential. There is another one where $!E$ is made of **sets** instead of multisets; it is not compatible with the differential structure.

Coh is canonically summable

- **Coh** has 0-morphisms: $0 = \emptyset \in \mathbf{Coh}(E, F)$.
- $l = 1$ & 1 so that $|l| = \{0, 1\}$ and $0 \subset_l 1$.
- The injections $\bar{\pi}_i = \{(*, i)\} \in \mathbf{Coh}(1, l)$ are jointly epic.

Remark

$s \in \mathbf{Coh}(l, E)$ is fully determined by the pair

$$s_0 = s \cdot \{0\}, s_1 = s \cdot \{1\} \in \mathbf{Cl}(E)$$

Moreover, since $0 \frown_E 1$ (which means $0 \supset_E 1$ and $0 \neq 1$) we have

$$s_0 \cup s_1 \in \text{Cl}(E) \quad \text{and} \quad s_0 \cap s_1 = \emptyset$$

Conversely if $x_0, x_1 \in \text{Cl}(E)$ satisfy $x_0 \cup x_1 \in \text{Cl}(E)$ and $x_0 \cap x_1 = \emptyset$ then

$$(\{0\} \times x_0) \cup (\{1\} \times x_1) \in \mathbf{Coh}(I, E)$$

Summability in Coh

We have seen that:

Fact

$x_0, x_1 \in Cl(E)$ are summable in E iff

$$x_0 \cup x_1 \in Cl(E) \quad \text{and} \quad x_0 \cap x_1 = \emptyset.$$

Remark

Each model of LL has its own notion of summability.

The CWA holds in **Coh**

Up to iso:

$$\text{Cl}(\mathbf{S}_I E) = \{(x_0, x_1) \in \text{Cl}(E)^2 \mid x_0 \cup x_1 \in \text{Cl}(E) \text{ and } x_0 \cap x_1 = \emptyset\}$$

Remark

Up to this iso, we have

$$\emptyset = (\emptyset, \emptyset)$$

$$(x_{00}, x_{01}) \cup (x_{10}, x_{11}) = (x_{00} \cup x_{10}, x_{01} \cup x_{11})$$

$$(x_{00}, x_{01}) \cap (x_{10}, x_{11}) = (x_{00} \cap x_{10}, x_{01} \cap x_{11}).$$

Summability in \mathbf{S}_1E

So $(x_{00}, x_{01}), (x_{10}, x_{11}) \in \text{Cl}(\mathbf{S}_1E)$ are summable in \mathbf{S}_1E if

$$(x_{00} \cup x_{10}, x_{10} \cup x_{11}) \in \text{Cl}(\mathbf{S}_1E)$$

$$(x_{00} \cap x_{10}, x_{10} \cap x_{11}) = (\emptyset, \emptyset)$$

That is

$$x_{00} \cup x_{10} \cup x_{10} \cup x_{11} \in \text{Cl}(E)$$

$$(x_{00} \cup x_{10}) \cap (x_{10} \cup x_{11}) = x_{00} \cap x_{10} = x_{10} \cap x_{11} = \emptyset$$

that is $(i, j) \neq (i', j') \Rightarrow x_{ij} \cap x_{i'j'} = \emptyset$.

Assume that

- $(x_{00}, x_{01}), (x_{10}, x_{11}) \in \text{Cl}(\mathbf{S}_I E)$ and
- $(x_{00} \cup x_{01}, x_{10} \cup x_{11}) \in \text{Cl}(\mathbf{S}_I E)$.

Then

- $x_{00} \cap x_{01} = x_{10} \cap x_{11} = \emptyset$
- $(x_{00} \cup x_{01}) \cap (x_{10} \cup x_{11}) = \emptyset$
- $x_{00} \cup x_{01} \cup x_{10} \cup x_{11} \in \text{Cl}(\mathbf{S}_I E)$

that is

- $x_{00} \cup x_{01} \cup x_{10} \cup x_{11} \in \text{Cl}(\mathbf{S}_I E)$
- $(i, j) \neq (i', j') \Rightarrow x_{ij} \cap x_{i'j'} = \emptyset$.

that is $(x_{00}, x_{01}), (x_{10}, x_{11}) \in \text{Cl}(\mathbf{S}_I E)$ are summable in $\mathbf{S}_I E$.

We already know that **Coh** has a unique differential structure wrt. $!$.

The commutative comonoid structure of $!$ is given by

$$\text{pr}_0 = \{(0, *)\} \in \mathbf{Coh}(!, 1)$$

$$\tilde{!} = \{(0, (0, 0)), (1, (1, 0)), (1, (0, 1))\} \in \mathbf{Coh}(!, ! \otimes !).$$

Remember it induces the monad structure $\zeta_E \in \mathbf{Coh}(E, \mathbf{S}_! E)$ and $\theta_E \in \mathbf{Coh}(\mathbf{S}_!^2 E, \mathbf{S}_! E)$.

As expected for $((x_{00}, x_{01}), (x_{10}, x_{11})) \in \text{Cl}(\mathbf{S}_!^2 E)$ we have

$$\theta \cdot ((x_{00}, x_{01}), (x_{10}, x_{11})) = (x_{00}, x_{10} + x_{01}) \in \text{Cl}(\mathbf{S}_! E)$$

The differential structure of **Coh**

We define $\delta \subseteq |I \multimap !!|$:

$$\delta = \{(0, n[0]) \mid n \in \mathbb{N}\} \cup \{(1, n[0] + [1]) \mid n \in \mathbb{N}\}$$

where $n[a] = \overbrace{[a, \dots, a]}^{n \times}$.

$\delta \in \mathbf{Coh}(I, !!)$ because

- $m \circ_{!!} m'$ for all $m, m' \in |!!|$
- and $n[0] \frown_{!E} n'[0] + [1]$ for all $n, n' \in \mathbb{N}$.

δ is a coalgebra

The main thing to check is

$$\begin{array}{ccc} | & \xrightarrow{\delta} & !! \\ \delta \downarrow & & \downarrow !\delta \\ !! & \xrightarrow{\text{dig}_!} & !!! \end{array}$$

that is, given $i \in \{0, 1\}$ and $M \in \mathcal{M}_{\text{fin}}(\mathcal{M}_{\text{fin}}(\{0, 1\}))$,

$$(i, M) \in !\delta \delta \Leftrightarrow (i, M) \in \text{dig}_! \delta$$

where

$$\text{dig}_E = \{(m, [m_1, \dots, m_k]) \in |E| \times !!E \mid m = m_1 + \dots + m_k\}.$$

main case

The main case is when $i = 1$.

$(1, M) \in !\delta\delta$ means $\exists k \in \mathbb{N}$ such that

$$(k[0] + [1], M) \in !\delta$$

that is:

$$M = [m_1, \dots, m_{k+1}] \text{ with } (0, m_i) \in \delta \text{ for } i = 1, \dots, k \\ \text{and } (1, m_{k+1}) \in \delta$$

that is: $\exists k \in \mathbb{N} \exists n_1, \dots, n_{k+1} \in \mathbb{N}$

$$M = [n_1[0], \dots, n_k[0], n_{k+1}[0] + [1]]$$

And $(1, M) \in \text{dig}_1 \delta$ means $\exists k \in \mathbb{N}$ such that

$$(k[0] + [1], M) \in \text{dig}_1$$

that is:

$$M = [m_1, \dots, m_l] \text{ with } m_1 + \dots + m_l = k[0] + [1]$$

that is: $\exists l \in \mathbb{N}^+ \exists n_1, \dots, n_l \in \mathbb{N}$

$$M = [n_1[0], \dots, n_{l-1}[0], n_l[0] + [1]]$$

The diagram commutes!

The induced differential dist. law

Remember that δ induces a distributive law

$\partial_E = \text{cur } u \in \mathbf{Coh}(!\mathbf{S}_1 E, \mathbf{S}_1 !E)$ where

$$u : !(I \multimap E) \otimes I \rightarrow !E$$

is

$$!(I \multimap E) \otimes I \xrightarrow{\text{Id} \otimes \delta} !(I \multimap E) \otimes !I \xrightarrow{\mu^2} !((I \multimap E) \otimes I) \xrightarrow{!ev} !E$$

Notice that $\mu_{E,F}^2 \in \mathbf{Coh}(!E \otimes !F, !(E \otimes F))$ is

$$\mu_{E,F}^2 = \{(((a_1, \dots, a_n), [b_1, \dots, b_n]), [(a_1, b_1), \dots, (a_n, b_n)]) \mid n \in \mathbb{N}, [a_1, \dots, a_n] \in !E \text{ and } [b_1, \dots, b_n] \in !F\}$$

$$\begin{aligned}
\partial_E = & \{ ((([(0, a_1), \dots, (0, a_k)], (0, [a_1, \dots, a_k])) \mid \\
& \quad k \in \mathbb{N} \text{ and } [a_1, \dots, a_k] \in |!E| \} \\
& \cup \{ ((([(0, a_1), \dots, (0, a_k), (1, a_{k+1})], (1, [a_1, \dots, a_{k+1}])) \mid \\
& \quad k \in \mathbb{N} \text{ and } [a_1, \dots, a_{k+1}] \in |!E| \text{ and } a_{k+1} \notin \{a_1, \dots, a_k\} \} \\
& \in \mathbf{Coh}(!(| \multimap E), | \multimap !E).
\end{aligned}$$

The Kleisli category $\mathbf{Coh}_!$

Object: those of \mathbf{Coh} and $\mathbf{Coh}_!(E, F) = \mathbf{Coh}(!E, F)$.

$s \in \mathbf{Coh}_!(E, F)$ induces a **stable function**

$$\hat{s} : \text{Cl}(E) \rightarrow \text{Cl}(F)$$

$$x \mapsto \{b \in |F| \mid \exists m \in \mathcal{M}_{\text{fin}}(x) (m, b) \in s\}$$

Remark

Different s 's can induce the same stable function: \hat{s} forgets about the multiplicities in multisets.

If $s_1 = \{([a], b)\}$ and $s_2 = \{([a, a], b)\}$ then $\hat{s}_1 = \hat{s}_2$.

Differentiation on $\mathbf{Coh}_!$

Given $t \in \mathbf{Coh}_!(E, F) = \mathbf{Coh}(!E, F)$, remember that

$$\mathbf{D}t = (\mathbf{S}t) \partial_E \in \mathbf{Coh}(!\mathbf{S}_!E = (I \multimap E), \mathbf{S}_!F = (I \multimap F)).$$

Notice that

$$\mathbf{S}t = \{((i, m), (i, b)) \mid i \in \{0, 1\} \text{ and } (m, b) \in t\}.$$

So for $t \in \mathcal{L}_!(E, F) = \mathbf{Coh}(!E \multimap F)$ we have

$$\begin{aligned} \mathbf{D}t = & \{([(0, a_1), \dots, (0, a_k)], (0, b)) \mid ([a_1, \dots, a_k], b) \in t\} \\ & \cup \{([(0, a_1), \dots, (0, a_k), (1, a_{k+1})], (1, b)) \mid \\ & \quad ([a_1, \dots, a_k, a_{k+1}], b) \in t \text{ and } a_{k+1} \notin \{a_1, \dots, a_k\}\} \\ & \in \mathcal{L}_!(\mathbf{S}_!E, \mathbf{S}_!F) = \mathbf{Coh}(!(I \multimap E) \multimap (I \multimap F)). \end{aligned}$$

The stable derivative

Remember that

$$\text{Cl}(\mathbf{S}_1 E) \simeq \{(x, u) \mid x \cup u \in \text{Cl}(E) \text{ and } x \cap u = \emptyset\}.$$

In that way we get the stable function

$$\begin{aligned} \widehat{\mathbf{D}}t &: \text{Cl}(\mathbf{S}_1 E) \rightarrow \text{Cl}(\mathbf{S}_1 F) \\ (x, u) &\mapsto (\widehat{t}(x), t'(x) \cdot u) \end{aligned}$$

$$t'(x) \cdot u = \{b \in |F| \mid \exists m \in \mathcal{M}_{\text{fin}}(x), a \in u (m + [a], b) \in t\}$$

Remark

In such an (m, a) we have $a \notin \text{supp}(m)$ since $\text{supp}(m) \subseteq x$, $a \in u$ and $x \cap u = \emptyset$.

Local coherence space

Given $x \in \text{Cl}(X)$, one defines a coherence space E_x by

- $|E_x| = \{b \in |E| \mid \forall a \in x \ a \frown_E b\}$
- $a \circ_{E_x} a'$ if $a \circ_E a'$.

Then for $t \in \mathbf{Coh}(E, F)$ we have

$$t'(x) \in \mathbf{Coh}(E_x, F_{\widehat{t}(x)})$$

Remark

There is a dependent type intuition: the type of $t'(x)$ depends on x .

However this point of view hardly reflects the stability of $t'(x)$ wrt. x .

Whereas the compound construction $\mathbf{D}t$ does in a very simple way.

Remark

If $t_i = \{(i[a], a)\}$ for $i = 1, 2$ we get

$$t'_1(\emptyset) \cdot \{a\} = \{a\}$$

$$t'_2(\emptyset) \cdot \{a\} = \emptyset$$

whereas $\widehat{t}_1 = \widehat{t}_2$. The derivative stable function $\widehat{\mathbf{D}}t$ is associated with t and not the stable function \widehat{t} .

In some sense this derivative “does not see multiplicities”. This is due to the **uniformity** of the exponential. NB: there are non-uniform coherence spaces. . .

Probabilistic Coherence Spaces

Probabilistic Coherence Spaces (PCS)

$$X = (|X|, PX)$$

- $|X|$ is a set (usually at most countable)
- $PX \subseteq (\mathbb{R}_{\geq 0})^{|X|}$
- $\forall a \in |X| \ 0 < \sup_{x \in PX} x_a < \infty$
- PX is \downarrow -closed (for the pointwise order)
- PX contains the (pointwise) lub of any increasing ω -sequence in PX
- $x, y \in PX$ and $\lambda \in [0, 1] \Rightarrow \lambda x + (1 - \lambda)y \in PX$

Morphisms

$X \multimap Y$ defined by:

- $|X \multimap Y| = |X| \times |Y|$
- and $t \in (\mathbb{R}_{\geq 0})^{|X| \times |Y|}$ is in $P(X \multimap Y)$ if

$$\forall x \in PX \quad t \cdot x \in PY$$

where $t \cdot x = (\sum_{a \in |X|} t_{a,b} x_a)_{b \in |Y|} \in (\mathbb{R}_{\geq 0})^{|Y|}$.

Fact

$X \multimap Y$ so defined is a PCS.

$$\mathbf{Pcoh}(X, Y) = P(X \multimap Y).$$

Notations

- If $s \in \mathbf{Pcoh}(X, Y)$ and $t \in \mathbf{Pcoh}(Y, Z)$ then $ts \in \mathbf{Pcoh}(X, Z)$ given by

$$(ts)_{a,c} = \sum_{b \in |Y|} s_{a,b} t_{b,c}$$

- $\text{Id}_X \in \mathbf{Pcoh}(X, X)$ is $(\delta_{a,a'})_{(a,a') \in |X \multimap X|}$.

This defines a category.

Cartesian product

- Terminal object \top such that $|\top| = \emptyset$.
- $|X_0 \& X_1| = \{0\} \times |X_0| \cup \{1\} \times |X_1|$ so that $(\mathbb{R}_{\geq 0})^{|X_0 \& X_1|} \simeq (\mathbb{R}_{\geq 0})^{|X_0|} \times (\mathbb{R}_{\geq 0})^{|X_1|}$
- $\text{pr}_i \in (\mathbb{R}_{\geq 0})^{|X_0 \& X_1| \times |X_i|}$ given by

$$(\text{pr}_i)_{(j,a),a'} = \delta_{i,j} \delta_{a,a'}$$

- $y \in (\mathbb{R}_{\geq 0})^{|X_0 \& X_1|}$ is in $P(X_0 \& X_1)$ if $\text{pr}_i \cdot y \in PX_i$ for $i = 0, 1$.
- If $t_i \in \mathbf{Pcoh}(Y, X_i)$ for $i = 0, 1$ then $\langle t_0, t_1 \rangle \in \mathbf{Pcoh}(Y, X_0 \& X_1)$ is given by $\langle t_0, t_1 \rangle_{b,(i,a)} = (t_i)_{a,b}$.

Remark

$$P(X_0 \& X_1) \simeq PX_0 \times PX_1$$

Up to this iso, the cartesian product is completely standard:

- $\text{pr}_i \cdot (x_0, x_1) = x_i$
- $\langle t_0, t_1 \rangle \cdot y = (t_0 \cdot y, t_1 \cdot y)$ for $t_i \in \mathbf{Pcoh}(Y, X_i)$

Tensor product

Given $x_i \in PX_i$ for $i = 0, 1$, let $x_0 \otimes x_1 \in (\mathbb{R}_{\geq 0})^{|X_0| \times |X_1|}$ given by

$$(x_0 \otimes x_1)_{(a_0, a_1)} = x_{0a_0} x_{1a_1}$$

- $|X_0 \otimes X_1| = |X_0| \times |X_1|$
- $P(X_0 \otimes X_1)$ minimal such that $x_0 \otimes x_1 \in P(X_0 \otimes X_1)$ for all $x_i \in PX_i$ for $i = 0, 1$.

Fact

$\mathbf{Pcoh}(Z \otimes X, Y) \simeq \mathbf{Pcoh}(Z, X \multimap Y)$.

The object I

$1 = (\{*\}, [0, 1])$. Notice that $P(1 \multimap X) \simeq PX$

$I = 1 \& 1$ so that $PI = [0, 1] \times [0, 1]$

$\bar{\pi}_0, \bar{\pi}_1 \in \mathbf{Pcoh}(1, I) \simeq PI$, actually $\bar{\pi}_0 = (1, 0)$ and $\bar{\pi}_1 = (0, 1)$.

Fact

$\bar{\pi}_0, \bar{\pi}_1$ are jointly monic:

by linearity, $t \in \mathbf{Pcoh}(I, X)$ is fully determined by

$t_0 = t \cdot (1, 0) \in PX$ and $t_1 = t \cdot (0, 1) \in PX$.

Moreover $t_0 + t_1 \in PX$ since $t_0 + t_1 = t \cdot (1, 1)$ since $(1, 1) \in PI$.

$P(I \multimap X) \simeq \{(x_0, x_1) \in PX \mid x_0 + x_1 \in PX\}$.

Canonical Witness Axiom

Up to this iso we have

$$(x_{00}, x_{01}) + (x_{10}, x_{11}) = (x_{00} + x_{10}, x_{01} + x_{11})$$

and so

$$\begin{aligned} ((x_{00}, x_{01}), (x_{10}, x_{11})) \in \mathbf{PS}_I^2 X &\Leftrightarrow (x_{00} + x_{10}, x_{01} + x_{11}) \in \mathbf{PS}_I X \\ &\Leftrightarrow x_{00} + x_{10} + x_{01} + x_{11} \in PX \\ &\Leftrightarrow (x_{00} + x_{01}, x_{10} + x_{11}) \in \mathbf{PS}_I X \end{aligned}$$

and hence the CWA holds.

The induced monad $\mathbf{S}_1 : \mathbf{Pcoh} \rightarrow \mathbf{Pcoh}$ given by $\mathbf{S}_1 X = (I \multimap X)$ behaves exactly as expected:

$$\zeta_X \in \mathbf{Pcoh}(X, \mathbf{S}_1 X)$$

$$\zeta_X \cdot x = (x, 0)$$

$$\theta_X \in \mathbf{Pcoh}(\mathbf{S}_1^2 X, \mathbf{S}_1 X) \quad \theta_X \cdot ((x_{00}, x_{01}), (x_{10}, x_{11})) = (x_{00}, x_{10} + x_{01})$$

The differentiation coalgebra

It is defined essentially as in **Coh**, and is a coalgebra for the same reason.

$$\delta \in (\mathbb{R}_{\geq 0})^{|\!-\circ\!|}$$

defined by

$$\delta_{i,m} = \begin{cases} 1 & \text{if } i = 0 \text{ and } \exists n \in \mathbb{N} \ m = n[0] \\ 1 & \text{if } i = 1 \text{ and } \exists n \in \mathbb{N} \ m = n[0] + [1] \\ 0 & \text{otherwise} \end{cases}$$

Remark (Surprise)

The case $i = 1$ implements the differential, so I expected to have sthg like n as coeff. instead of 1. But it's not the case!

The exponential functor

- $!X = \mathcal{M}_{\text{fin}}(|X|)$ (no uniformity restriction).
- If $x \in PX$ and $m \in !X$ then $x^m = \prod_{a \in |X|} x_a^{m(a)} \in \mathbb{R}_{\geq 0}$
- $x^! = (x^m)_{m \in !X}$
- and $P(!X)$ is minimal such that $\forall x \in PX \ x^! \in P(!X)$.

Given $t \in \mathbf{Pcoh}(X, Y)$ we need $!t \in \mathbf{Pcoh}(!X, !Y)$ such that

$$\forall x \in PX \quad !t \cdot x^! = (t \cdot x)^!$$

Fact

This determines fully $!t$.

Simple computations give, for $t \in \mathbf{Pcoh}(X, Y) \subseteq (\mathbb{R}_{\geq 0})^{|X| \times |Y|}$ and $(m, p) \in |!X \multimap !Y| = \mathcal{M}_{\text{fin}}(!X) \times \mathcal{M}_{\text{fin}}(!Y)$:

$$(!t)_{m,p} = \sum_{r \in L(m,p)} \begin{bmatrix} p \\ r \end{bmatrix} t^r$$

where

$$L(m, p) = \{r \in \mathcal{M}_{\text{fin}}(|X| \times |Y|) \mid \sum_{b \in |Y|} r(a, b) = m(a) \text{ and } \sum_{a \in |X|} r(a, b) = p(b)\}$$

and

$$\begin{bmatrix} p \\ r \end{bmatrix} = \prod_{b \in |Y|} \frac{p(b)!}{\prod_{a \in |X|} r(a, b)!} \in \mathbb{N}$$

The evaluation morphism

$$\text{ev} \in \mathbf{Pcoh}((I \multimap X) \otimes I, X) \quad \text{ev}_{((i,a),j),b} = \delta_{a,b} \delta_{i,j}$$

Then $! \text{ev}_{M,m} \neq 0$ implies

$$M = [((0, a_1), 0), \dots, ((0, a_k), 0), ((1, b_1), 1), \dots, ((1, b_n), 1)] = (l, r)$$

$$m = [a_1, \dots, a_k, b_1, \dots, b_n] = l + r$$

Setting $l = [a_1, \dots, a_k]$ and $r = [b_1, \dots, b_n]$. We have

$$! \text{ev}_{M,m} = \binom{l+r}{l} = \prod_{a \in |X|} \binom{l(a)+r(a)}{l(a)}$$

The differential functor

Remember that $\partial_X = \text{cur } f \in \mathcal{L}(!\mathbf{S}_I X, \mathbf{S}_I !X)$ where f is

$$!(I \multimap X) \otimes I \xrightarrow{\text{Id} \otimes \delta} !(I \multimap X) \otimes !I \xrightarrow{\mu^2} !((I \multimap X) \otimes I) \xrightarrow{! \text{ev}} !X$$

Using the above computation of $! \text{ev}$ and definition of δ we get

$$(\partial_X)_{(l,r),(i,m)} = \begin{cases} 1 & \text{if } i = 0, r = [], m = l \\ m(a) & \text{if } i = 1, r = [a], m = l + [a] \\ 0 & \text{otherwise.} \end{cases}$$

A $t \in \mathbf{Pcoh}_!(X, Y) = P(!X \multimap Y)$ is completely characterized by the associated **analytic function**

$$\hat{t} : PX \rightarrow PY$$

$$x \mapsto t \cdot x^! = \left(\sum_{m \in |!X|} t_{m,b} x^m \right)_{b \in |Y|}$$

Then $\mathbf{D}t \in \mathbf{Pcoh}_!(\mathbf{S}_!X, \mathbf{S}_!Y)$ is characterized by the analytic function (setting $f = \hat{t} : PX \rightarrow PY$)

$$\begin{aligned} \widehat{\mathbf{D}t} : P(\mathbf{S}_!X) &\rightarrow P(\mathbf{S}_!Y) \\ (x, u) &\mapsto (f(x), f'(x) \cdot u) \end{aligned}$$

where

$$f'(x) \cdot u = \left(\sum_{a \in |X|} \left(\sum_{l \in |X|} (m(a) + 1) t_{m+[a], bX^m} \right) u_a \right)_{b \in |Y|}$$

is just the standard differential of \hat{t} .

Differential as a linear map

Given $x \in PX$ we define X_x , the local PCS at x :

$$|X_x| = \{a \in |X| \mid \exists \varepsilon > 0 \ x + \varepsilon e_a \in PX\}$$

$$PX_x = \{u \in (\mathbb{R}_{\geq 0})^{|X_x|} \mid x + u \in PX\}$$

and then $f'(x) \in \mathbf{Pcoh}(X_x, Y_{f(x)})$ satisfies (for $b \in |Y_{f(x)}|$)

$$(f'(x) \cdot u)_b = \left(\frac{d}{dt} f(x + tu)_b \right)_{t=0} \quad \text{standard Gateaux derivative.}$$

The fact that $\mathbf{Dt} \in \mathbf{Pcoh}_!(\mathbf{S}_1 X, \mathbf{S}_1 Y)$ also tells us that this derivative is analytic in x .

Strong similarity with Tangent Categories

Mfd: category of smooth manifolds and smooth maps.

There is a *tangent bundle* functor

$$\mathbf{T} : \mathbf{Mfd} \rightarrow \mathbf{Mfd}$$

$$X \mapsto \{(x, u) \mid x \in X \text{ and } u \in \mathbf{T}_x X\}$$

$\mathbf{T}_x X$ = tangent space at x to X . A **vector space**.

And if $f \in \mathbf{Mfd}(X, Y)$,

$$\mathbf{T}f : \mathbf{T}X \rightarrow \mathbf{T}Y$$

$$(x, u) \mapsto (f(x), f'(x) \cdot u)$$

Looks very much like our **D** functor!

Discrepancies

Is **T** a special case of **D**?

Of course not: given $(x, u) \in \mathbf{TX}$, it makes no sense to consider u alone (no 2nd projection $\mathbf{TX} \rightarrow X$) nor to compute $x + u \in X$ in general.

Is **D** a special case of **T**?

No, because our “tangent spaces” are only partial commutative monoids whereas $\mathbf{T}_x X$ is crucially a commutative monoid.

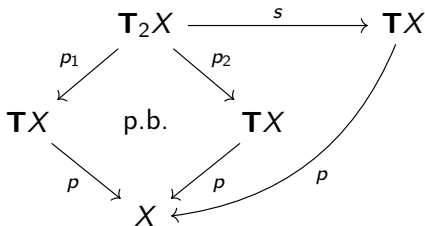
Remark

More philosophically, our approach is based on **S** acting on a “linear” category (a category of algebraic objects, the linear category of a model of LL).

This is typically not the case in the tangent bundle case.

More precisely, in *tangent categories* (= categorical axiomatization of the tangent bundle functor) we have a natural transformation $p_X : \mathbf{T}X \rightarrow X$, intuitively $p_X(x, u) = x$.

It is required that there is a pull-back and an *addition morphism* s



This s is a **total** addition operation in the fibers of p .